MO417 – Complexidade de Algoritmos Segundo Semestre de 2011 Nona Lista de Exercícios

Caminhos Mínimos

- 1. Give a simple example of a directed graph with negative-weight edges for which Dijkstra's algorithm produces incorrect answers.
- 2. We are given a directed graph G = (V, E) on which each edge $(u, v) \in E$ has an associated value r(u, v), which is a real number in the range $0 \le r(u, v) \le 1$ that represents the reliability of a communication channel from vertex u to vertex v. We interpret r(u, v) as the probability that the channel from u to v will not fail, and we assume that these probabilities are independent. Give an efficient algorithm to find the most reliable path between two given vertices.
- 3. Suppose we change the line:

"while $Q \neq \emptyset$ "

of Dijkstra's algorithm to the following:

"while |Q| > 1"

This change causes the while loop to execute |V| - 1 times instead of |V| times. Is this proposed algorithm correct?

- 4. Let G = (V, E) be a weighted, directed graph with weight function $w : E \to \{0, 1, ..., W\}$ for some nonnegative integer W. Modify Dijkstra's algorithm to compute the shortest paths from a given source vertex s in O(WV + E) time.
- 5. Modify your algorithm from previous exercise to run in $O((V + E) \lg W)$ time.
- 6. Modify the Bellman-Ford algorithm so that it sets d[v] to $-\infty$ for all vertices v for which there is a negative-weight cycle on some path from the source to v.
- 7. Given a weighted, directed graph G = (V, E) with no negative-weight cycles, let m be the maximum over all pairs of vertices $u, v \in V$ of the minimum number of edges in a shortest path from u to v. Here, the shortest path is by weight, not the number of edges. Suggest a simple change to the Bellman-Ford algorithm that allows it to terminate in m + 1 passes.
- 8. Give an efficient algorithm to count the total number of paths in a directed acyclic graph. Analyze your algorithm.