



# Predicate Logic

## Quantifiers $\forall, \exists$



## Predicates

Predicates are  
**Propositions with variables**

Example:

$$P(x,y) \stackrel{\text{::=}}{=} x + 2 = y$$

“is defined to be”



## Predicates

$$P(x, y) ::= [x + 2 = y]$$

$x = 1$  and  $y = 3$ :  $P(1,3)$  is true

$x = 1$  and  $y = 4$ :  $P(1,4)$  is false  
 $\neg P(1,4)$  is true



## Quantifiers

$\forall x$  For ALL  $x$

$\exists y$  There EXISTS some  $y$



## Quantifiers

$x, y$  range over **Domain of Discourse**

$$\forall x \exists y. x < y$$

<u>Domain</u>	<u>Truth value</u>
integers $\mathbb{Z}$	True
positive integers $\mathbb{Z}^+$	True
negative integers $\mathbb{Z}^-$	False
negative reals $\mathbb{R}^-$	True



$\forall \alpha$  versus  $\exists d$

~~$\forall \alpha \in \text{attack}$~~   ~~$\exists d \in \text{defense}$~~ .  
 $d$  protects against  $\alpha$

For every attack, I have a defense:  
against MYDOOM, use Defender  
against ILOVEYOU, use Norton  
against BABLAS, use Zonealarm ...

$\forall \exists$  is expensive!



$\exists \forall$

$\exists d \in \text{defense } \forall a \in \text{attack.}$   
 $d$  protects against  $a$

I have *one* defense good  
against every attack.

Example:  $d$  is MITviruscan,  
protects against *all* viruses

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So  $\exists \forall$  is better here

$\exists d \in \text{defense } \forall a \in \text{attack.}$   
 $d$  protects against  $a$

I have *one* defense good  
against every attack.

That's what we want!

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Math vs. English

Poet:

$\overbrace{G} \quad \overbrace{Au}$   
"All that glitters is not gold."

$\forall x. G(x) \rightarrow Au(x)$

No!: gold glitters like gold

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Math vs. English

Poet:

$\text{necessarily}$   
"All that glitters is not gold."

$\neg [\forall x. G(x) \rightarrow Au(x)]$

(Poetic license)

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Math vs. English

Poet: "There is a season for every  
purpose under heaven"

$\exists s \in \text{Season } \forall p \in \text{Purpose. } s \text{ is for } p$

So some season, say Spring, is good for  
all Purposes?

NO, Spring is no good for snow shoveling

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Poetic license again:

Poet: "There is a season for every  
purpose under heaven"

$\exists s \in \text{Season } \forall p \in \text{Purpose. } s \text{ is for } p$

Poet's meaning flips the quantifiers

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4	9	13	7
12		10	5
3	1	6	14
15	8	11	2

## Poetic license again:

**Poet:** “There is a season for every purpose under heaven”

$\forall p \in \text{Purpose} \exists s \in \text{Season}. s$  is for  $p$   
 for snow shoveling, *Winter* is good  
 for planting, *Spring* is good  
 for leaf watching, *Fall* is good  
 etc.

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4	9	13	7
12		10	5
3	1	6	14
15	8	11	2

## Team Problems

# Problems 1 & 2

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4	9	13	7
12		10	5
3	1	6	14
15	8	11	2

## Propositional Validity

$$(A \rightarrow B) \vee (B \rightarrow A)$$

**True** *no matter what* the truth values of  $A$  and  $B$  are

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4	9	13	7
12		10	5
3	1	6	14
15	8	11	2

## Predicate Calculus Validity

$$\forall z [Q(z) \wedge P(z)] \rightarrow [\forall x.Q(x) \wedge \forall y.P(y)]$$

**True** *no matter what*

- the Domain is,
- or the predicates are.

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4	9	13	7
12		10	5
3	1	6	14
15	8	11	2

## Not Valid

$$\forall z [Q(z) \vee P(z)] \rightarrow [\forall x.Q(x) \vee \forall y.P(y)]$$

**Proof:** Give *countermodel*, where  
 $\forall z [Q(z) \vee P(z)]$  is **true**,  
 but  $\forall x.Q(x) \vee \forall y.P(y)$  is **false**.

Namely, let domain  $::= \{e, \pi\}$ ,  
 $Q(z) ::= [z = e]$ ,  
 $P(z) ::= [z = \pi]$ .

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4	9	13	7
12		10	5
3	1	6	14
15	8	11	2

## Predicate Calculus Validity

$$\forall z [Q(z) \wedge P(z)] \rightarrow [\forall x.Q(x) \wedge \forall y.P(y)]$$

**Proof strategy:** We assume

$$\forall z [Q(z) \wedge P(z)]$$

to prove

$$\forall x.Q(x) \wedge \forall y.P(y)$$

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4	9	13	7
12	10	6	1
3	1	8	14
15	5	11	2

## Universal Generalization (UG)

$$\frac{A \rightarrow R(c)}{A \rightarrow \forall x.R(x)}$$

providing  $c$  does not occur in  $A$

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4	9	13	7
12	10	6	1
3	1	8	14
15	5	11	2

## Validities

$$\forall z [Q(z) \wedge P(z)] \rightarrow [\forall x.Q(x) \wedge \forall y.P(y)]$$

*Proof:* Assume  $\forall z [Q(z) \wedge P(z)]$ .

So  $Q(z) \wedge P(z)$  holds for all  $z$  in the domain.

Now let  $c$  be some domain element. So

$Q(c) \wedge P(c)$  holds, and therefore  $Q(c)$  by itself holds.

But  $c$  could have been any element of the domain.

So we conclude  $\forall x.Q(x)$ . (by **UG**)

We conclude  $\forall y.P(y)$  similarly. Therefore,

$$\forall x.Q(x) \wedge \forall y.P(y) \quad \text{QED.}$$

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4	9	13	7
12	10	6	1
3	1	8	14
15	5	11	2

## More Validities

$$\forall x[P(x) \vee A] \leftrightarrow [\forall x.P(x)] \vee A$$

providing  $x$  does not occur in  $A$

$$[\neg \forall x.P(x)] \leftrightarrow [\exists x.\neg P(x)]$$

(version of DeMorgan)

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4	9	13	7
12	10	6	1
3	1	8	14
15	5	11	2

## Team Problems

# Problems

# 4 & 3

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