

CENTER GENOME

WITH RESPECT TO THE

RANK DISTANCE

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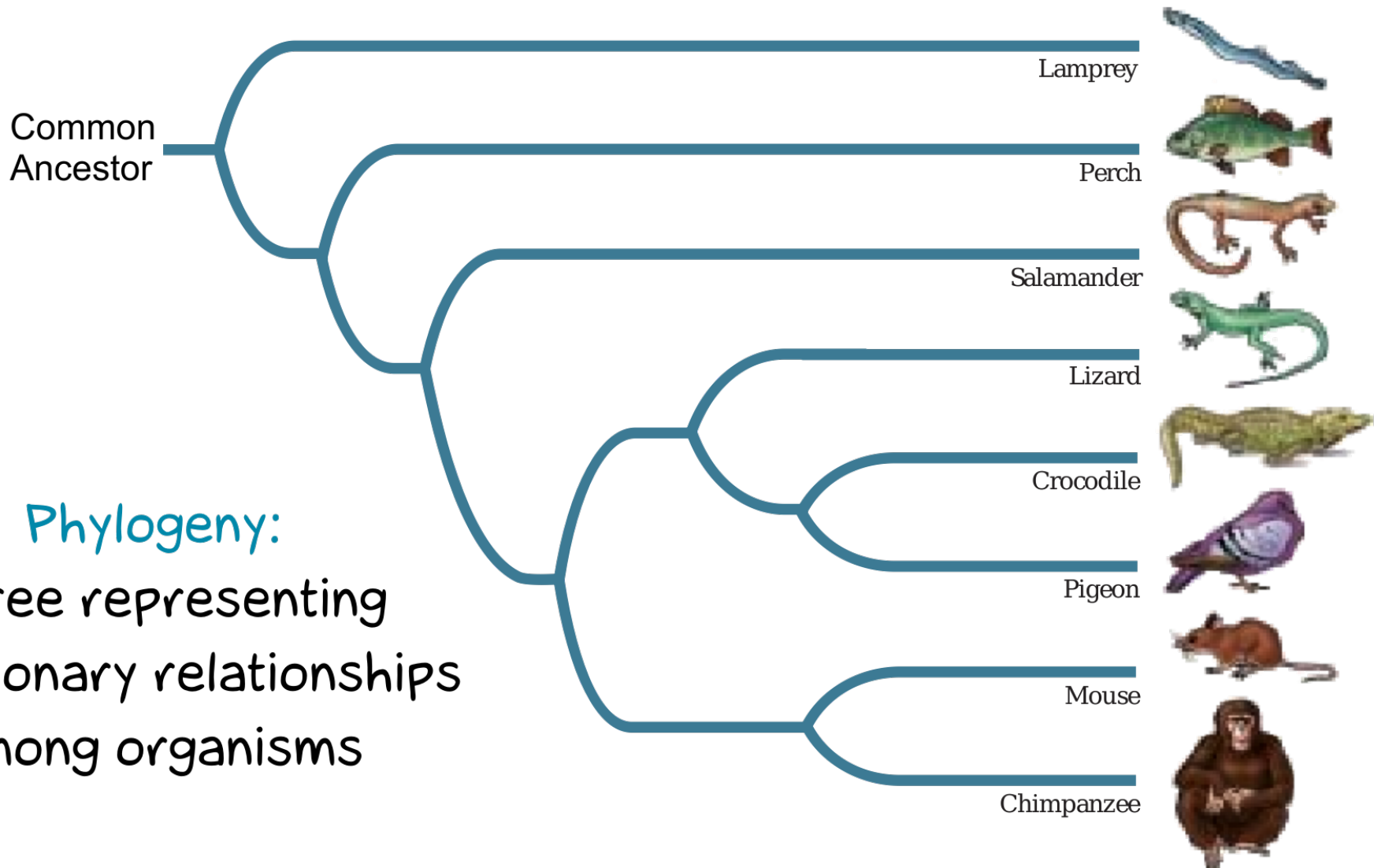


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How to infer ancestral genomes?

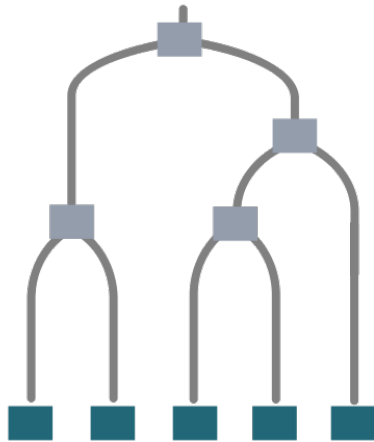


Phylogeny:
a tree representing
evolutionary relationships
among organisms

Internal nodes:
ancestors (usually extinct)

Leaves:
recent species (known genomes)

How to infer ancestral genomes?

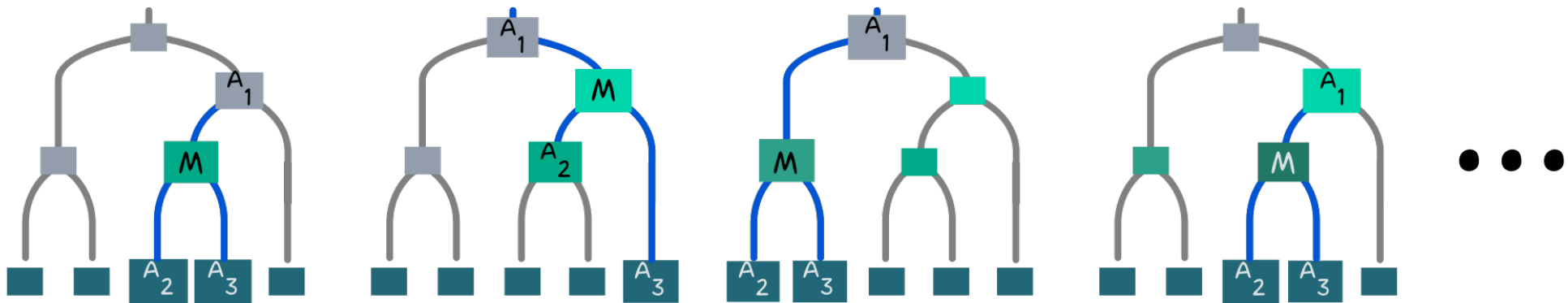


■ Known genomes

■ Genomes to be inferred (initially arbitrary)

Usual way to infer ancestors:

Repeatedly compute the median genome M until convergence is reached



Genome M minimizes the sum of the evolutionary distances:

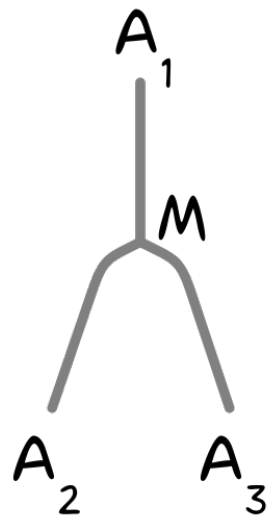
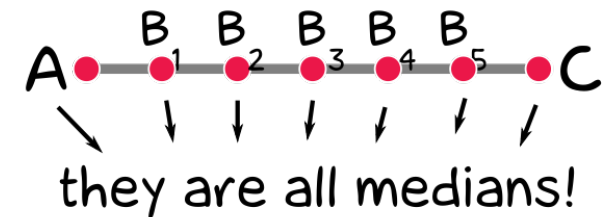
$$d(A_1, M) + d(A_2, M) + d(A_3, M)$$

Are median genomes a good way to infer ancestors?

2 input genomes: solutions are not relevant

Any genome in an optimal sorting scenario minimizes

$d(A, B) + d(B, C)$, including the input genomes



3 or more input genomes: hard
NP-hard for most rearrangement
distances (reversal, DCJ, etc.)

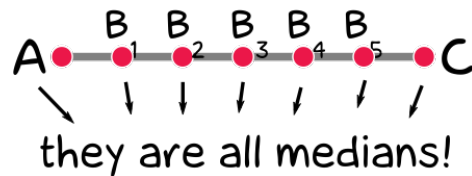
Center genome: an alternative to the median

Median genome

Input: genomes A_1, A_2, \dots

Goal: find a genome M that minimizes
 $d(A_1, M) + d(A_2, M) + d(A_3, M) + \dots$

2 input
genomes:



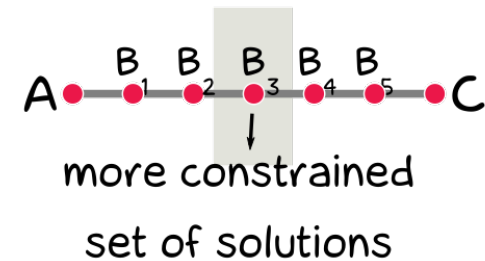
NP-hard

3 or more
input genomes:

Center genome

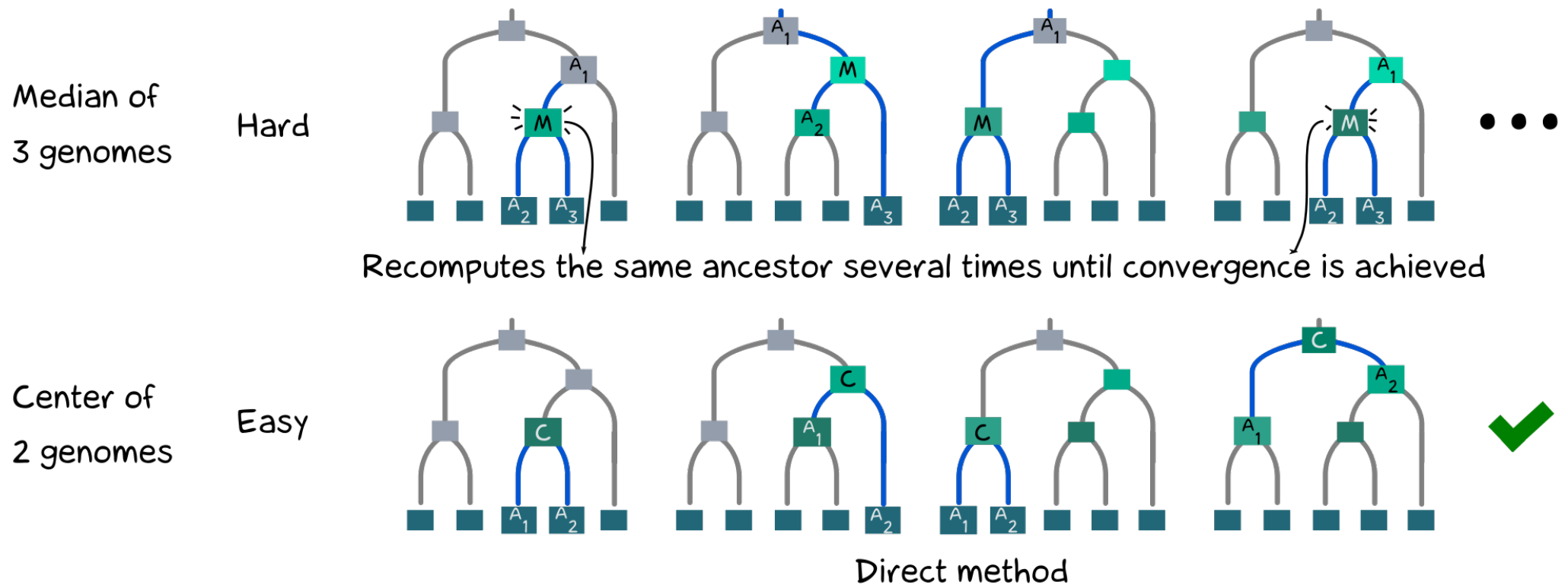
genomes A_1, A_2, \dots

find a genome M that minimizes
 $\max(d(A_1, M), d(A_2, M), d(A_3, M), \dots)$



Open (NP-hard?)

Ancestral inference: center genomes are an appealing alternative to the median



Center genomes with respect to the rank distance



Gabidulin Ernst

Rank distance:

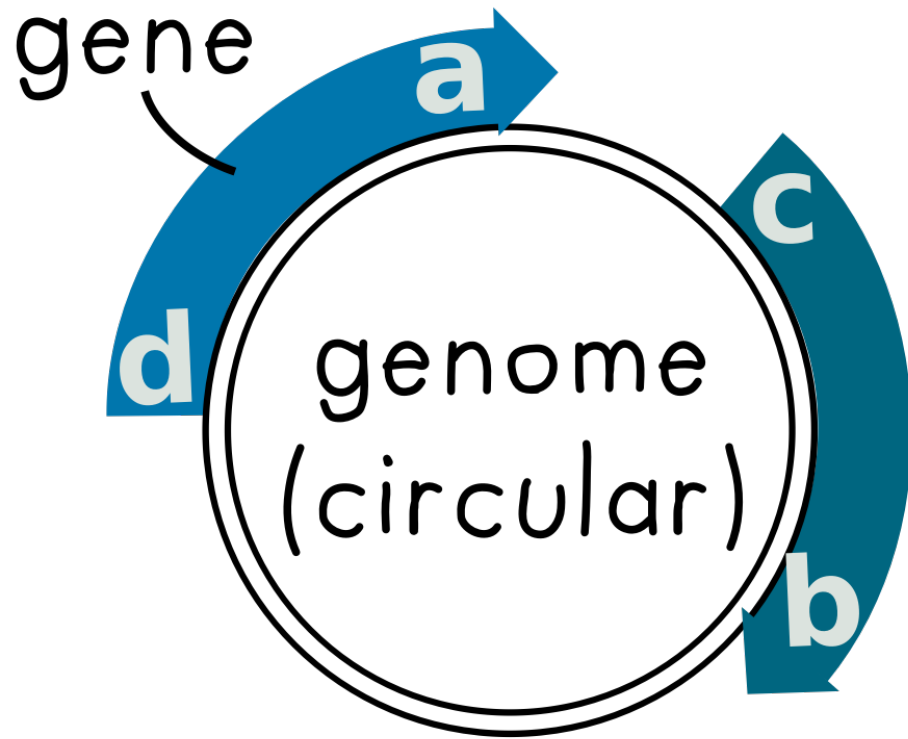
very successfully used in coding theory since at least 1985

$$d(A, B) = \text{rank}(A - B)$$

↓ ↓
matrices

$\text{rank}(X)$: dimension of the image (or column space) of X

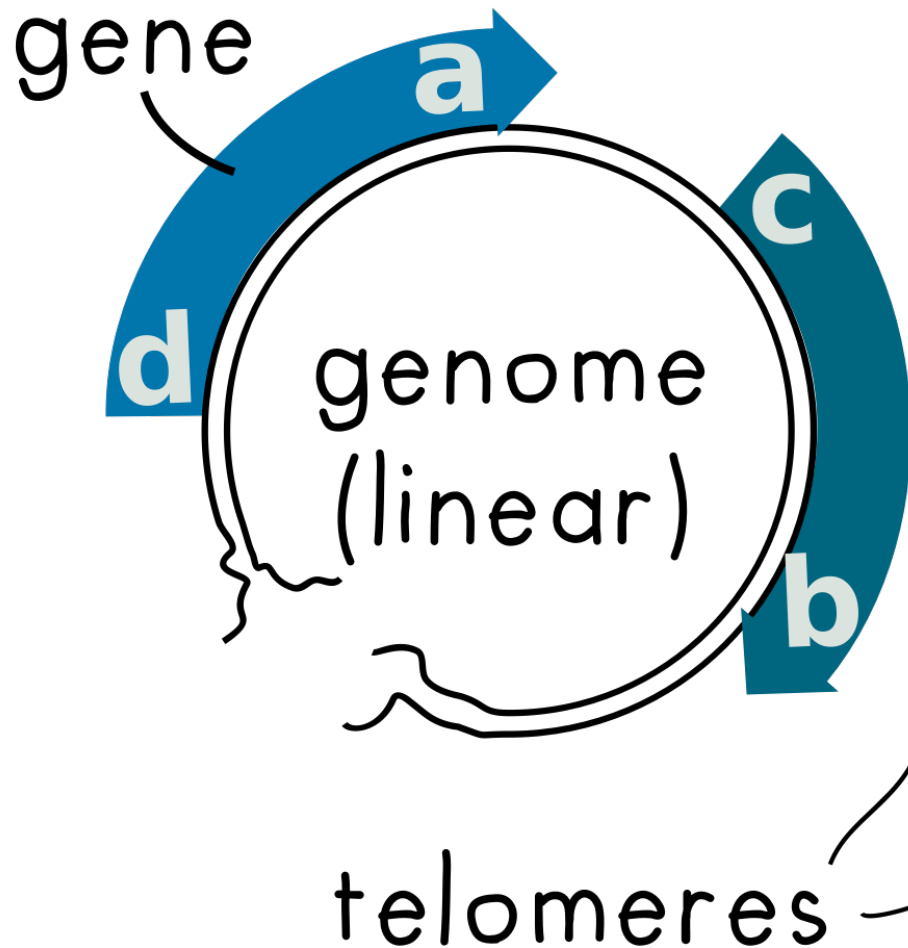
How to represent genomes as matrices?



	a	b	c	d
a	0	0	1	0
b	0	0	0	1
c	1	0	0	0
d	0	1	0	0

gene extremities

How to represent genomes as matrices?



	a	b	c	d
a	0	0	1	0
b	0	1	0	0
c	1	0	0	0
d	0	0	0	1

How to compute the rank distance of two genomes?

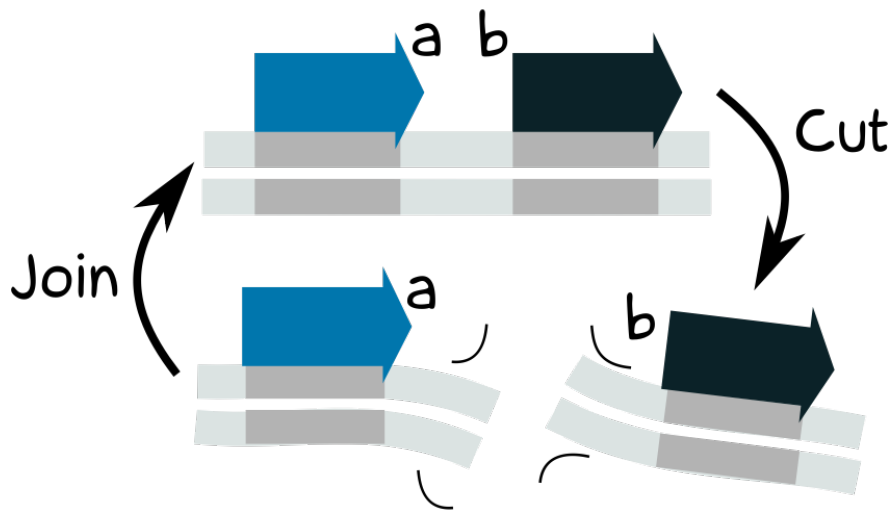


$$d(M_1, M_2) = \text{rank} \begin{pmatrix} & \mathbf{a} & \mathbf{b} & \mathbf{c} & \mathbf{d} \\ \mathbf{a} & 0 & 0 & \mathbf{1} & 0 \\ \mathbf{b} & 0 & 0 & 0 & \mathbf{1} \\ \mathbf{c} & \mathbf{1} & 0 & 0 & 0 \\ \mathbf{d} & 0 & \mathbf{1} & 0 & 0 \end{pmatrix} - \begin{pmatrix} & \mathbf{a} & \mathbf{b} & \mathbf{c} & \mathbf{d} \\ \mathbf{a} & 0 & 0 & \mathbf{1} & 0 \\ \mathbf{b} & 0 & \mathbf{1} & 0 & 0 \\ \mathbf{c} & \mathbf{1} & 0 & 0 & 0 \\ \mathbf{d} & 0 & 0 & 0 & \mathbf{1} \end{pmatrix}$$

$$= \text{rank} \begin{pmatrix} & \mathbf{a} & \mathbf{b} & \mathbf{c} & \mathbf{d} \\ \mathbf{a} & 0 & 0 & 0 & 0 \\ \mathbf{b} & 0 & \mathbf{-1} & 0 & \mathbf{1} \\ \mathbf{c} & 0 & 0 & 0 & 0 \\ \mathbf{d} & 0 & \mathbf{1} & 0 & \mathbf{-1} \end{pmatrix} = \mathbf{1}$$

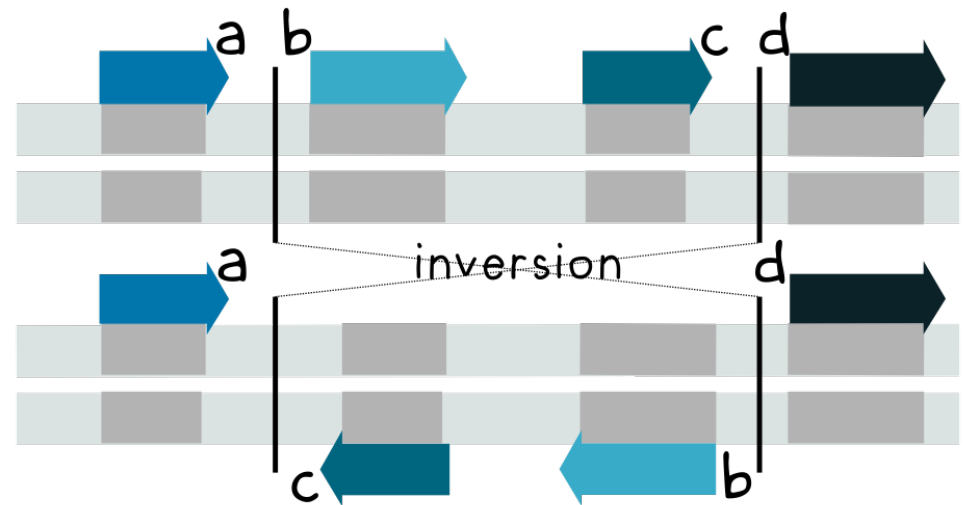
Applications of the rank distance to genome evolution

Rank distance can also be defined as the minimum number of cuts, joins, and double swaps, with weights 1, 1, and 2



$(a, b) \longleftrightarrow (a) \text{ and } (b)$

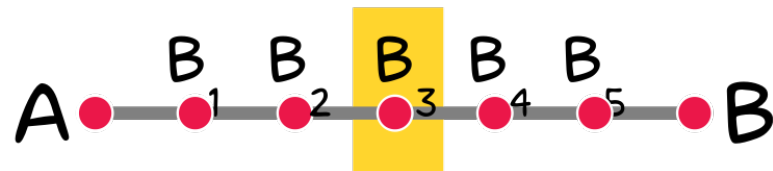
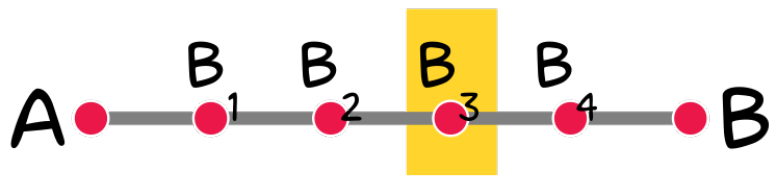
Double swap



$(a, b) \text{ and } (c, d) \begin{cases} \nearrow (a, c) \text{ and } (b, d) \\ \searrow (a, d) \text{ and } (c, b) \end{cases}$

Center conjecture:

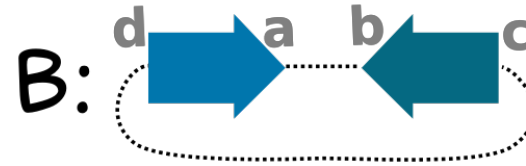
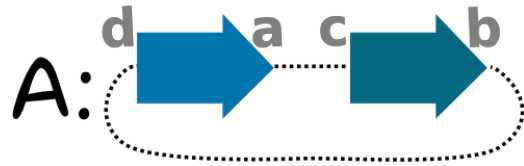
Is there always a genome exactly in the middle?



$$\max(d(A,C), d(B,C)) = \left\lceil \frac{d(A, B)}{2} \right\rceil$$

Center conjecture (counterexample):

Is there always a genome exactly in the middle?



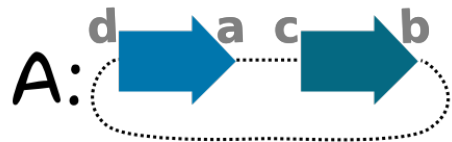
	a	b	c	d
a	0	0	1	0
b	0	0	0	1
c	1	0	0	0
d	0	1	0	0

	a	b	c	d
a	0	1	0	0
b	1	0	0	0
c	0	0	0	1
d	0	0	1	0

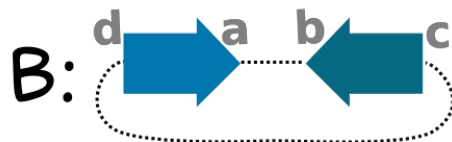
$$d(A, B) = \text{rank} \begin{pmatrix} & a & b & c & d \\ a & 0 & -1 & 1 & 0 \\ b & -1 & 0 & 0 & 1 \\ c & 1 & 0 & 0 & -1 \\ d & 0 & 1 & -1 & 0 \end{pmatrix} = 2$$

Center conjecture (counterexample):

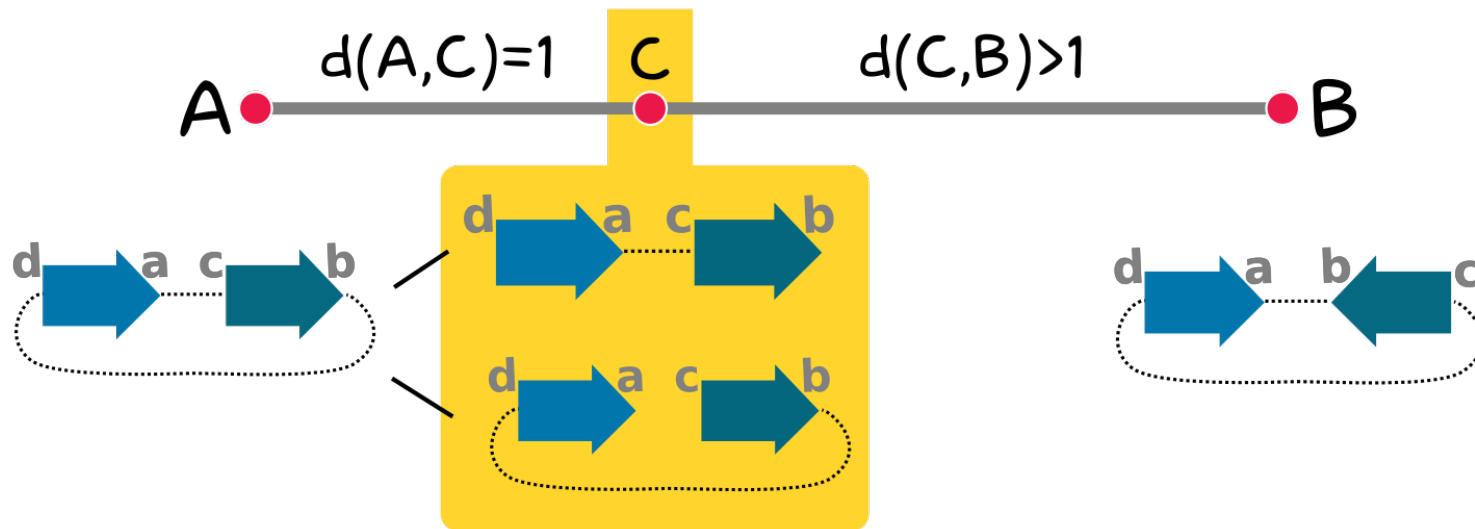
Is there always a genome exactly in the middle?



$$d(A, B) = \text{rank}(A-B) = 2$$



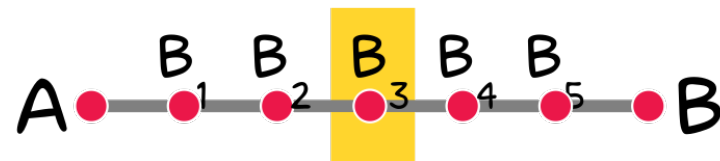
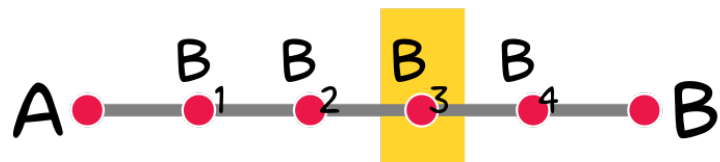
There is no genome with distance 1 from both A and B:



In this example: ~~$\max(d(A,C), d(B,C)) = \left\lfloor \frac{d(A,B)}{2} \right\rfloor$~~

Our goal:

Determine in which cases the center genome is exactly in the middle and, when not in the middle, how far it will be from it.



$$\max(d(A,C), d(B,C)) = \left\lceil \frac{d(A, B)}{2} \right\rceil$$

Our goal: Where is the center genome?

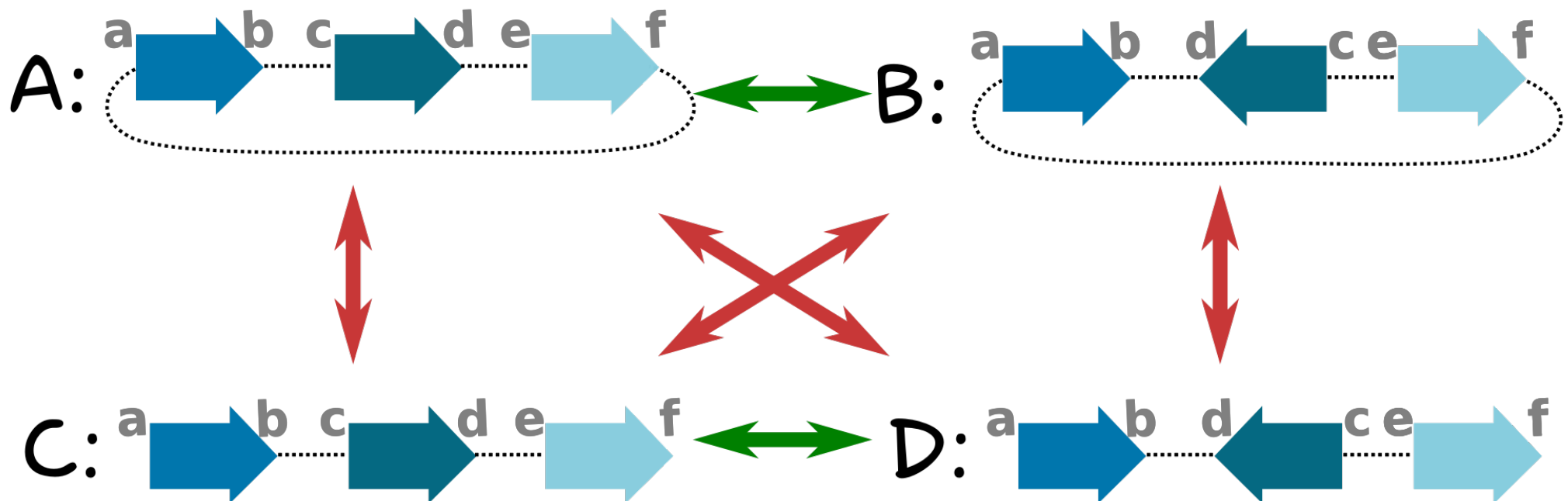
Given a pair of genomes, there are two cases to consider:

Co-tailed genomes

exactly the same telomeres

Not co-tailed genomes

different telomeres



Our goal: Where is the center genome?

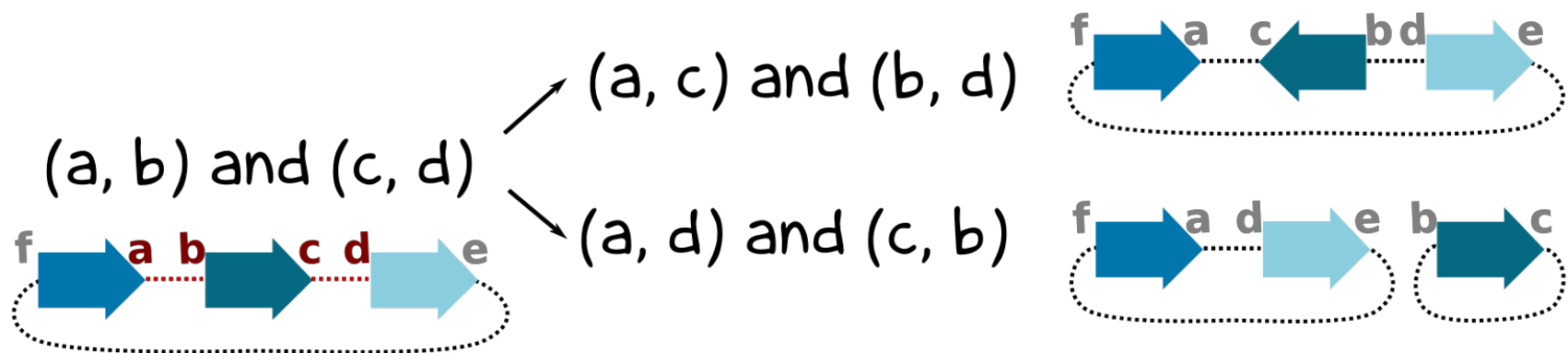


Co-tailed case (same telomeres)

Only double swaps occur in optimal sorting scenarios of co-tailed genomes!

Leonid Chindelevitch et al. (2018)

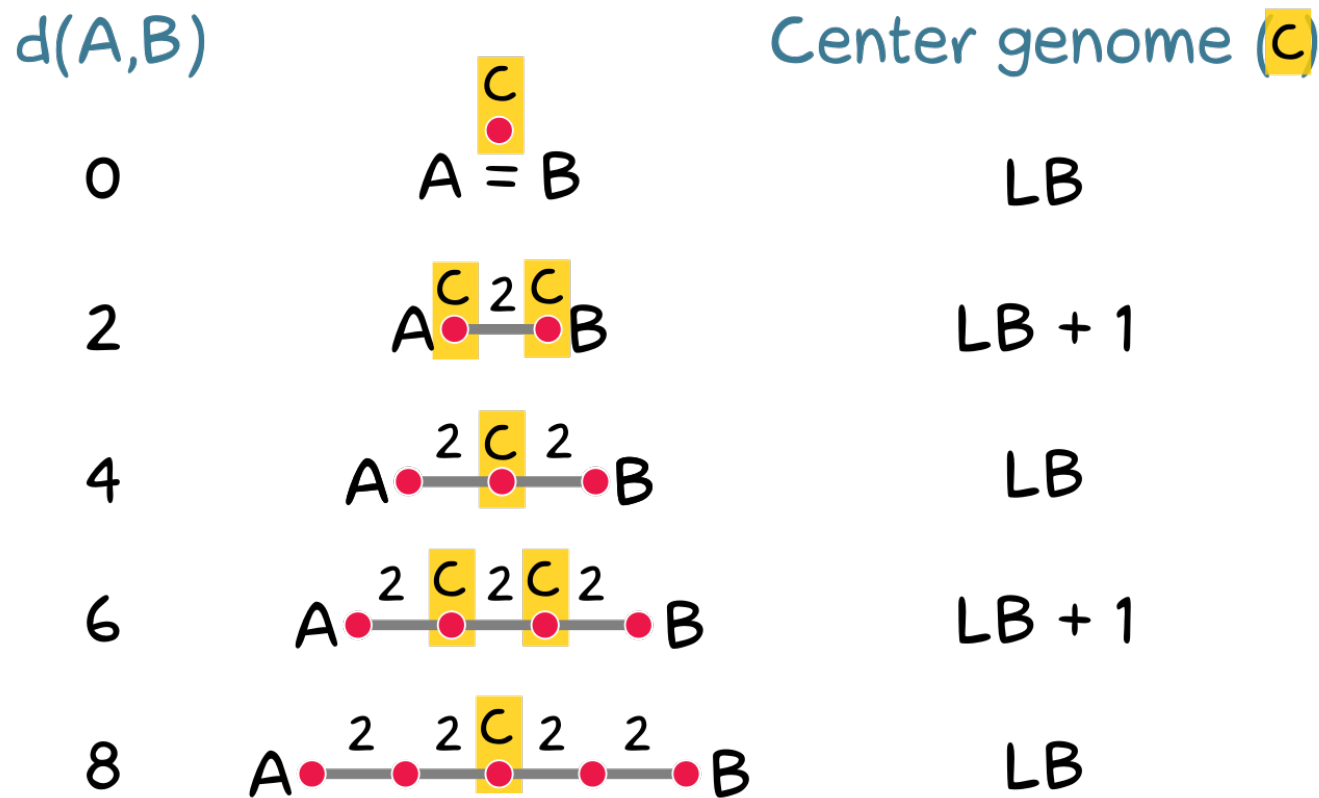
Double swap (weight = 2)



Our goal: Where is the center genome?

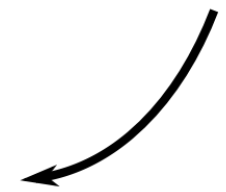
Co-tailed case (same telomeres)

As a double swap has weight 2, the distance is even!



Lower bound (LB):

$$\max(d(A,C), d(B,C)) = \left\lceil \frac{d(A,B)}{2} \right\rceil$$



Our goal: Where is the center genome?

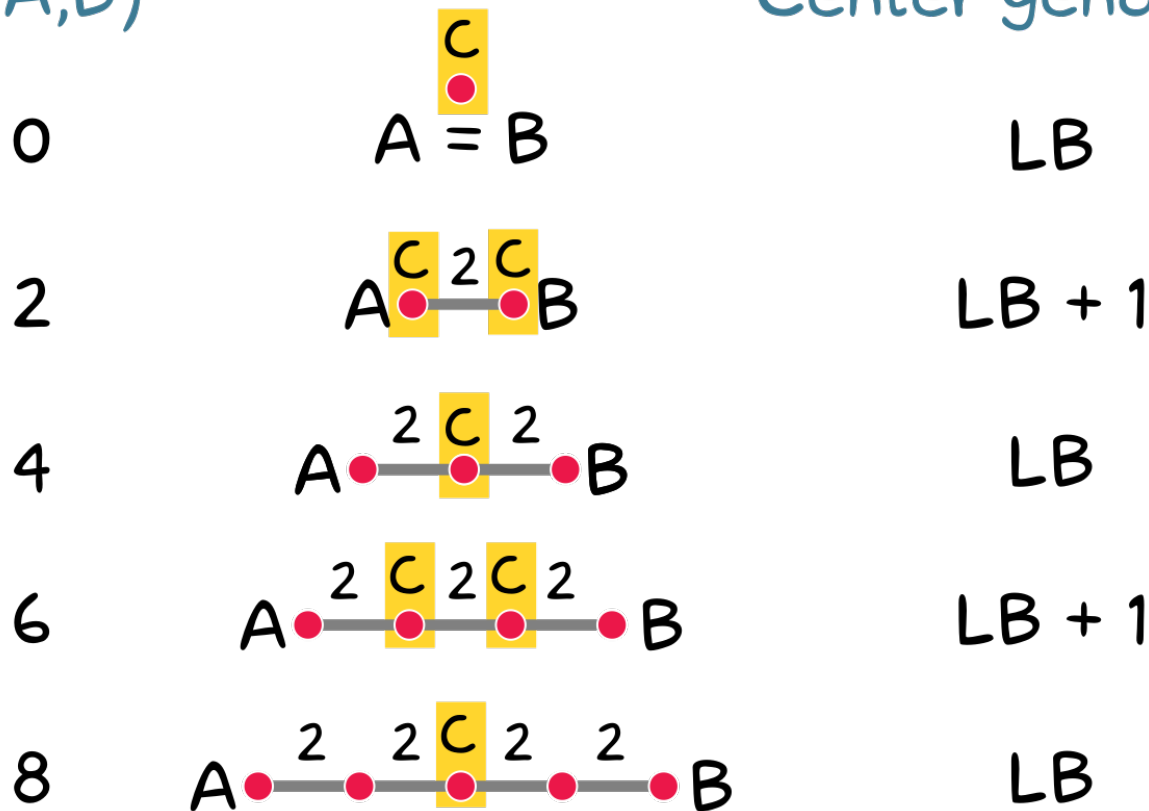
Co-tailed case (same telomeres)

$d(A,B)$ is a multiple of 4 : center genome reaches LB!

Otherwise: $LB + 1$

$d(A,B)$

Center genome (C)

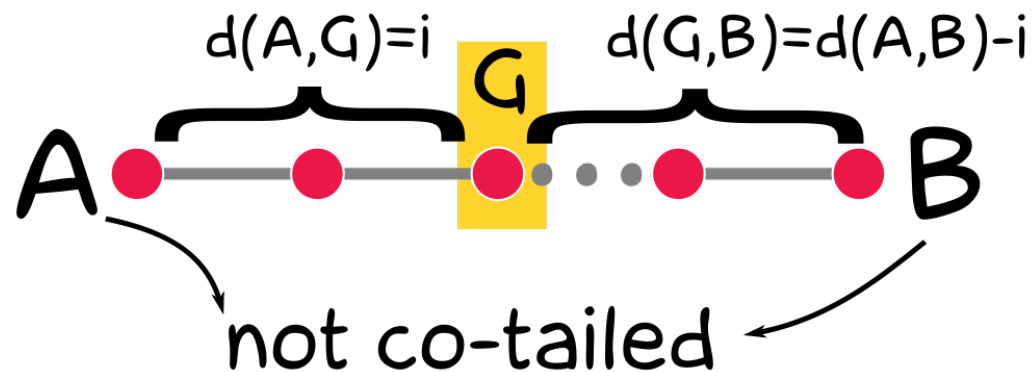


Our goal: Where is the center genome?

Not co-tailed case (different telomeres)

At every step $0 \leq i \leq d(A,B)$ there is a genome **G** such that

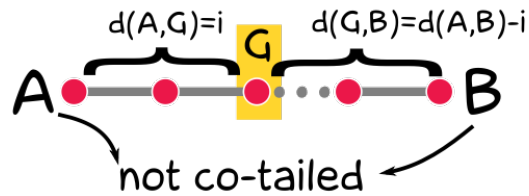
$$d(A, \mathbf{G})=i \text{ and } d(\mathbf{G}, B)=d(A,B)-i \dots$$



...so there is a center genome that reaches the LB!

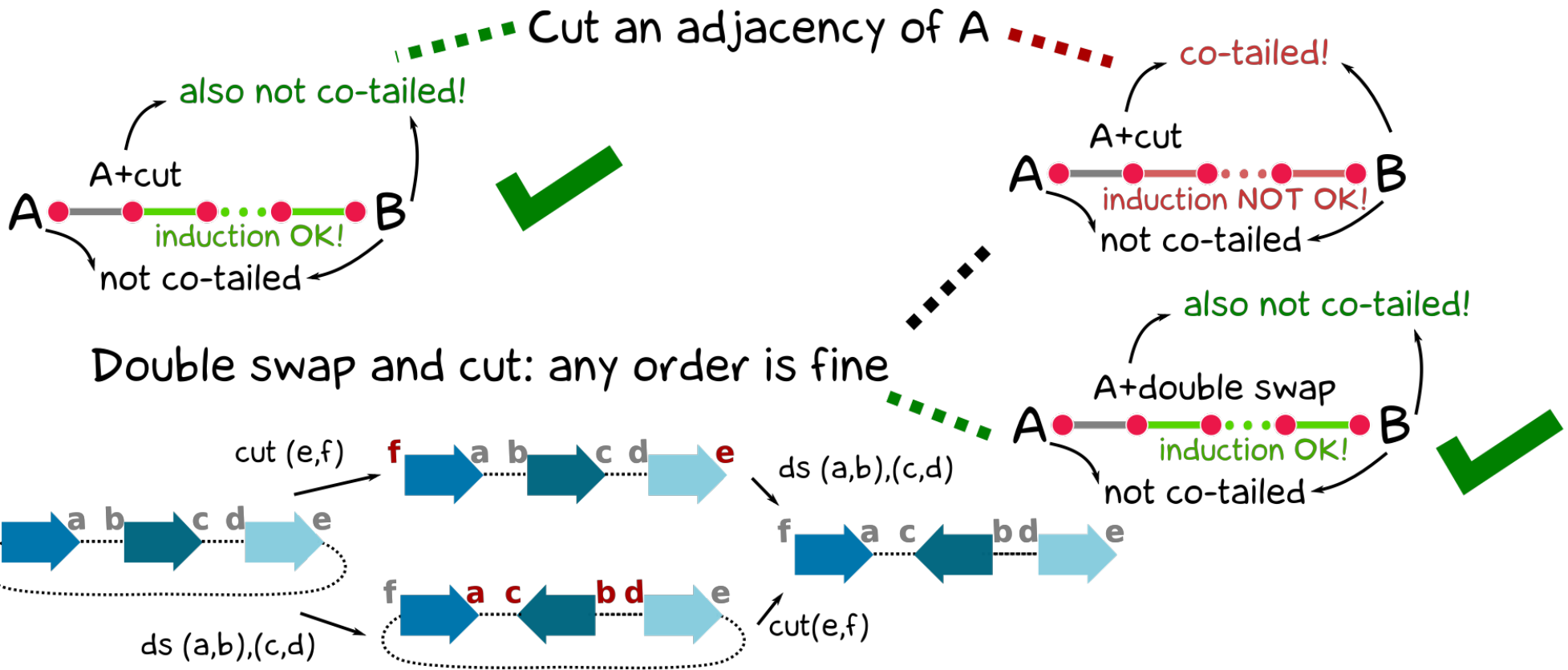
Our goal: Where is the center genome?

Not co-tailed case (different telomeres)



If A and B are not co-tailed, then there is a genome at every step between 0 and $d(A,B)$

Proof by induction (idea): Not co-tailed \rightarrow It needs a cut!



Summary: Where is the center genome?

Center genome (C)

(1) Co-tailed case (same telomeres)

$A \overset{2}{\bullet} \overset{2}{\text{C}} \overset{2}{\bullet} B$ $d(A,B)$ is a multiple of 4: LB

$A \overset{2}{\bullet} \overset{2}{\text{C}} \overset{2}{\text{C}} \overset{2}{\bullet} B$ $d(A,B)$ is not a multiple of 4: $LB+1$

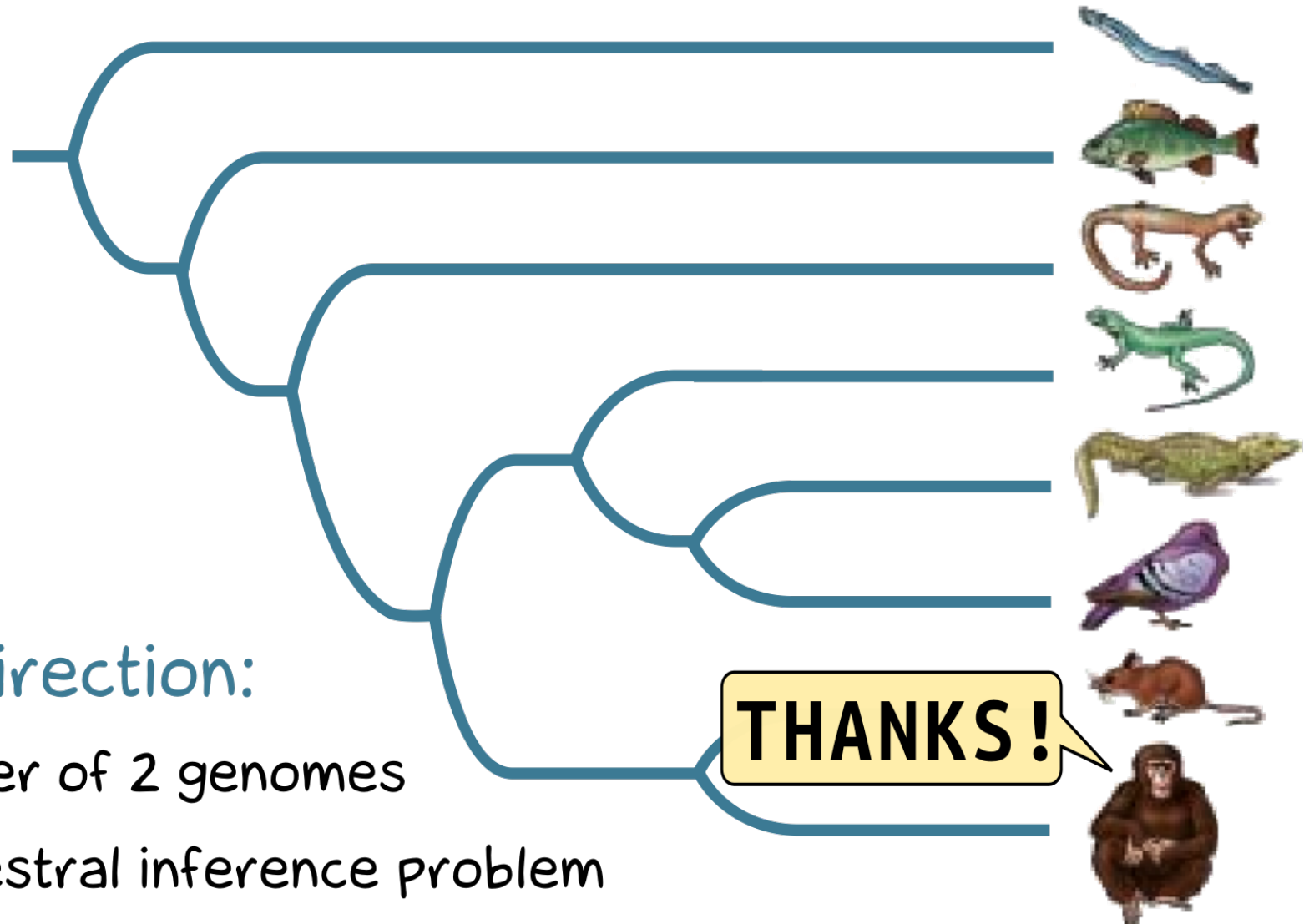
(2) Not co-tailed case (different telomeres) LB

Lower bound (LB):

$$\max(d(A,C), d(B,C)) = \left\lceil \frac{d(A,B)}{2} \right\rceil$$



Center genome with respect to the rank distance



Future direction:

Apply center of 2 genomes

to the ancestral inference problem