

MC504/MC514 - Sistemas Operacionais

Sistemas de Arquivos

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Sumário

- 1 Introdução
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- 3 Diretórios
- 4 Implementação

Sistemas de Arquivos

- Grande quantidade de informação
- Dados persistentes (não-voláteis)
- Acesso concorrente

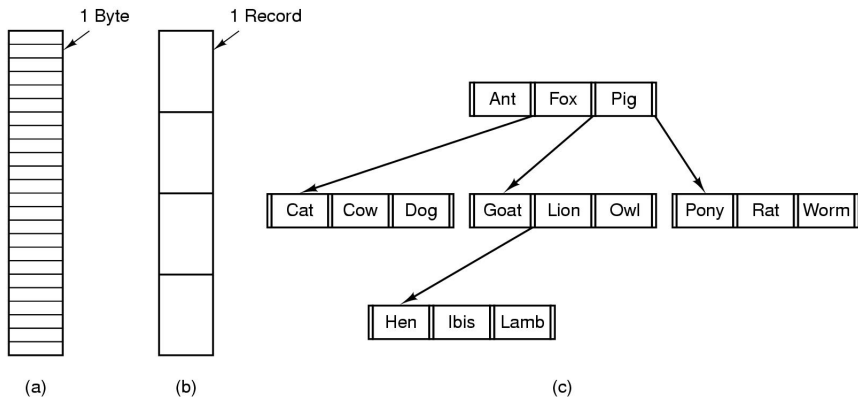
Nomes e extensões

Extension	Meaning
file.bak	Backup file
file.c	C source program
file.gif	Compuserve Graphical Interchange Format image
file.hlp	Help file
file.html	World Wide Web HyperText Markup Language document
file.jpg	Still picture encoded with the JPEG standard
file.mp3	Music encoded in MPEG layer 3 audio format
file.mpg	Movie encoded with the MPEG standard
file.o	Object file (compiler output, not yet linked)
file.pdf	Portable Document Format file
file.ps	PostScript file
file.tex	Input for the TEX formatting program
file.txt	General text file
file.zip	Compressed archive

Tanenbaum: Figura 6.1

Arquivos podem ter mais de uma extensão: `file.ps.gz`
Comando `file` verifica o tipo dos arquivos

Estruturas de arquivos



Tanenbaum: Figura 6.2

Tipos de arquivos

- regular
- diretório
- caracter
 - terminais, impressoras e rede
- bloco
 - discos

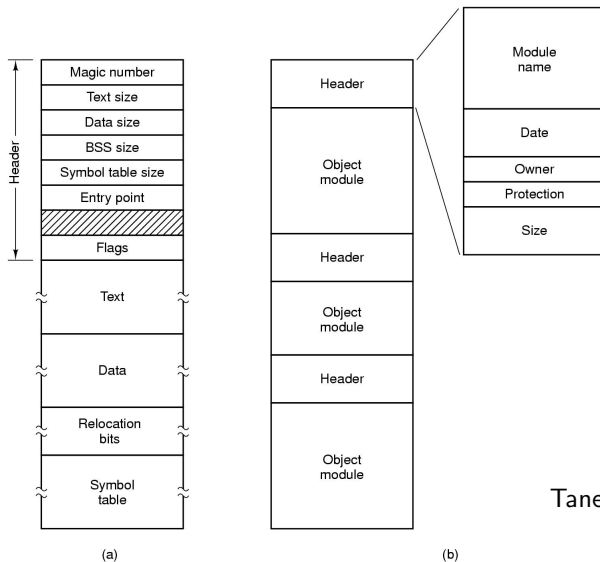
Use o comando `stat`:

```
$ stat arquivos.pdf
```

```
$ stat .
```

```
$ stat /dev/tty0
```

Exemplos: executável e archive



Tanenbaum: Figura 6.2

Acesso a arquivos

- Sequencial
 - Lê todos os bytes a partir do início
 - Fitas magnéticas
- Aleatório
 - Bytes podem ser lidos em qualquer ordem
 - Bancos de dados

Atributos de arquivos

Attribute	Meaning
Protection	Who can access the file and in what way
Password	Password needed to access the file
Creator	ID of the person who created the file
Owner	Current owner
Read-only flag	0 for read/write; 1 for read only
Hidden flag	0 for normal; 1 for do not display in listings
System flag	0 for normal files; 1 for system file
Archive flag	0 for has been backed up; 1 for needs to be backed up
ASCII/binary flag	0 for ASCII file; 1 for binary file
Random access flag	0 for sequential access only; 1 for random access
Temporary flag	0 for normal; 1 for delete file on process exit
Lock flags	0 for unlocked; nonzero for locked
Record length	Number of bytes in a record
Key position	Offset of the key within each record
Key length	Number of bytes in the key field
Creation time	Date and time the file was created
Time of last access	Date and time the file was last accessed
Time of last change	Date and time the file has last changed
Current size	Number of bytes in the file
Maximum size	Number of bytes the file may grow to

Tanenbaum: Figura 6.4

Veja os comandos `stat` e `make`

Operações sobre arquivos

- create
- delete
- open
- close
- read
- write
- append
- seek
- get attributes
- set attributes
- rename

Veja `linux-3.X.Y/include/linux/fs.h`

Programa copy

```
#define BUF_SIZE 4096
#define OUTPUT_MODE 0700

int main(int argc, char *argv[]) {
    int in_fd, out_fd, rd_count, wt_count;
    char buffer[BUF_SIZE];

    if (argc!=3) exit(1);

    in_fd = open(argv[1], O_RDONLY);
    if (in_fd < 0) exit(2);

    out_fd = creat(argv[2], OUTPUT_MODE);
    if (out_fd < 0) exit(3);
```

Programa copy

```
while((rd_count = read(in_fd, buffer, BUF_SIZE)) > 0) {  
    wt_count = write(out_fd, buffer, rd_count);  
    if (wt_count <= 0) exit(4);  
}
```

```
close(in_fd);  
close(out_fd);
```

```
if (rd_count == 0) exit(0);  
else exit(5);  
}
```

Streams and File Descriptors

File descriptors provide a primitive, low-level interface to input and output operations. [...]

The main advantage of using the stream interface is that the set of functions for performing actual input and output operations (as opposed to control operations) on streams is much richer and more powerful than the corresponding facilities for file descriptors. The file descriptor interface provides only simple functions for transferring blocks of characters, but the stream interface also provides powerful formatted input and output functions (`printf` and `scanf`) as well as functions for character- and line-oriented input and output.

Fonte: http://www.gnu.org/software/libc/manual/html_node/Streams-and-File-Descriptors.html

Streams

```
int fprintf(FILE *stream, const char *format, ...);
```

```
int fscanf(FILE *stream, const char *format, ...);
```

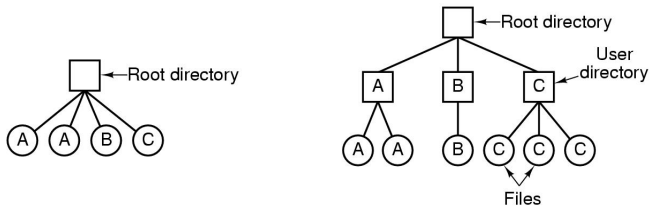
```
FILE *fopen(const char *path, const char *mode);
```

```
int fclose(FILE *stream);
```

Veja os exemplos `fscanf.c` e `fscanf2.c`

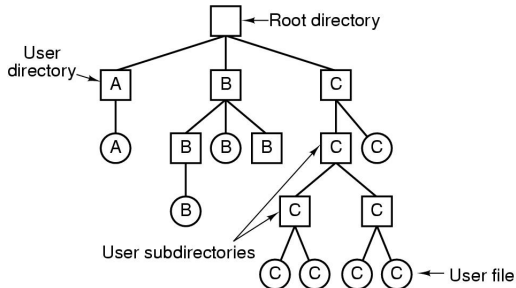
Diretórios

Nível único e dois níveis



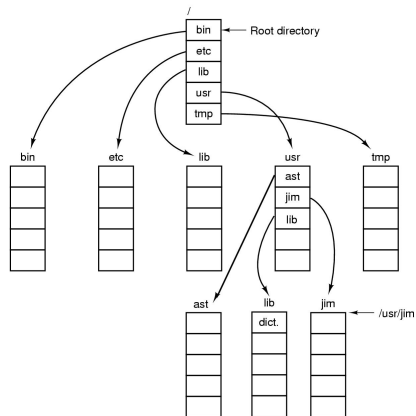
Tanenbaum: Figuras 6.7 e 6.8

Estrutura hierárquica



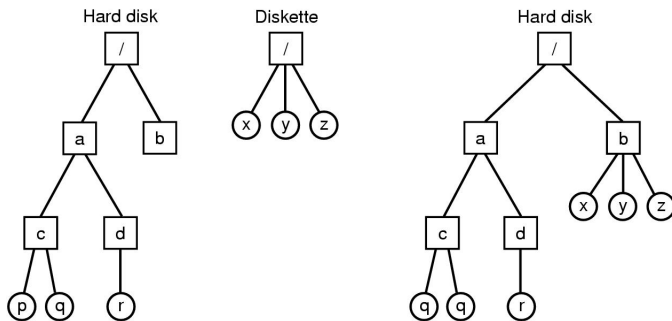
Tanenbaum: Figura 6.9

Caminhos



Tanenbaum: Figura 6.10

Mount



Tanenbaum: Figura 10.26

Mount

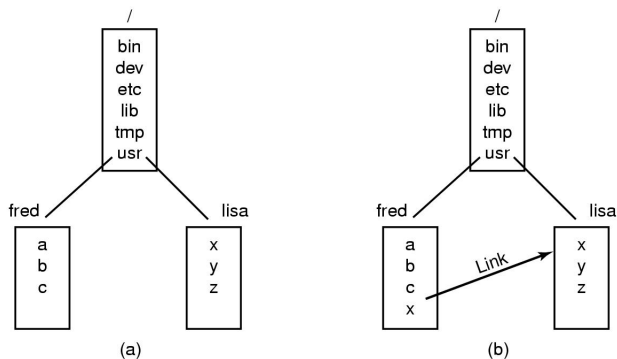
```
$ dd if=/dev/zero of=hd.dmp bs=1k count=4  
$ mk2fs hd.dmp  
$ mkdir -p mnt  
$ sudo mount -t ext2 -o loop hd.dmp mnt  
[sudo] password for islene:
```

Operações sobre diretórios

- create
- delete
- opendir
- closedir
- readdir
- rename
- link
- unlink

Veja o código `dir.c`

Links



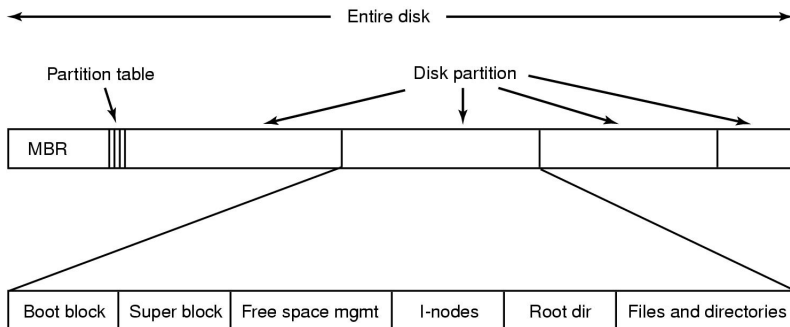
Tanenbaum: Figura 10.25

- links simbólicos ou hard links?
- caminhos absolutos ou relativos?
- como copiar?

Questões de Implementação

- Como os arquivos são armazenados
- Como o espaço livre é gerenciado
- Eficiência
- Confiabilidade

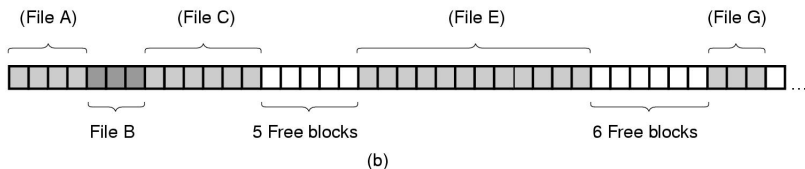
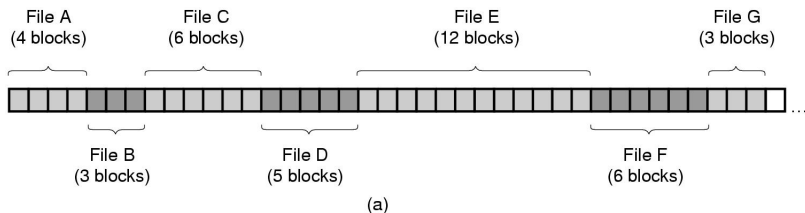
Layout



Tanenbaum: Figura 6.11

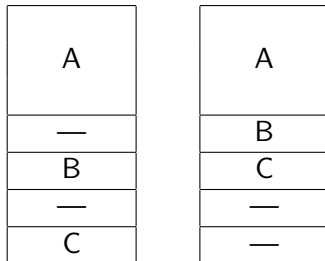
- MBR (Master Boot Record)
- Tabela de partições
- Boot block

Alocação contínua

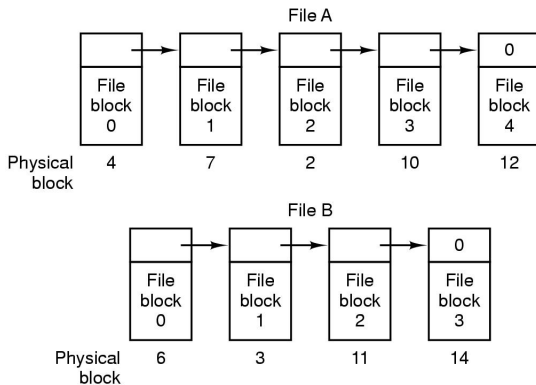


Tanenbaum: Figura 6.12

Compactação do disco

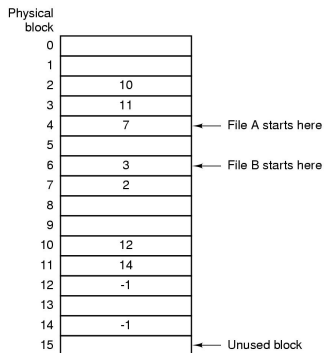


Lista ligada de blocos



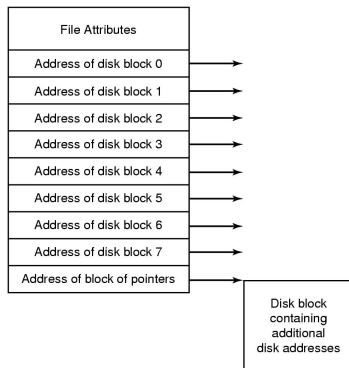
Tanenbaum: Figura 6.13

File Allocation Table (FAT)



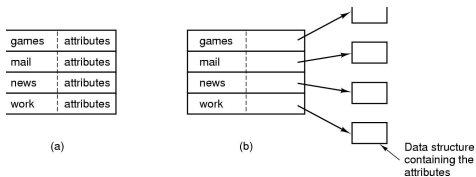
Tanenbaum: Figura 6.14

I-node



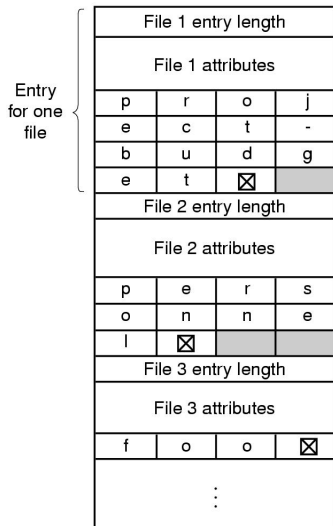
Tanenbaum: Figura 6.15

Implementação de diretórios

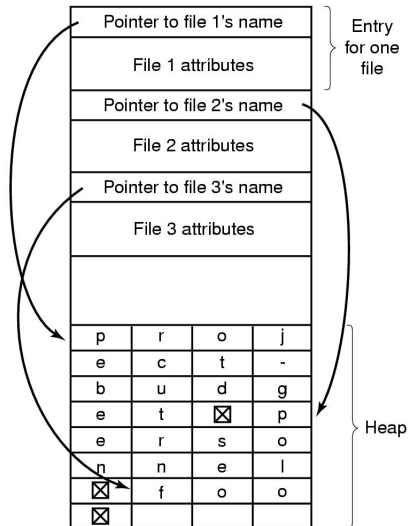


Tanenbaum: Figura 6.16

Nomes de tamanho variável

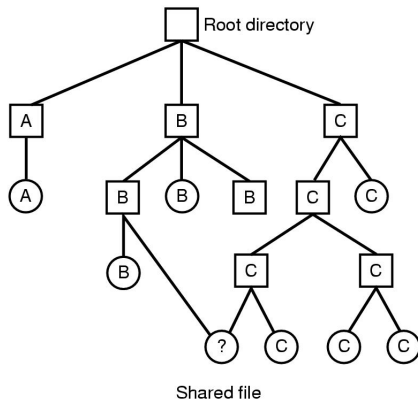


(a)



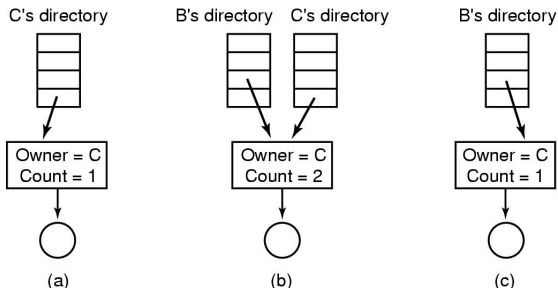
(b)

Arquivos compartilhados



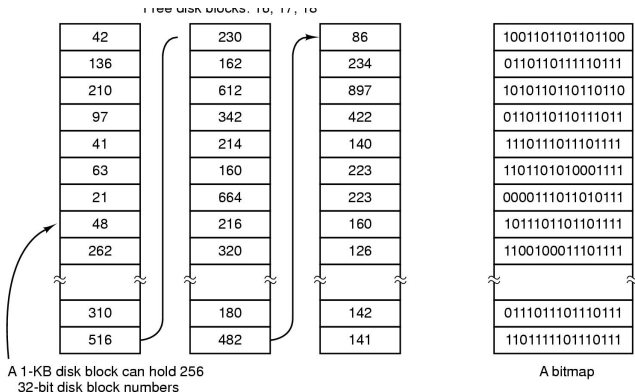
Tanenbaum: Figura 6.18

Arquivos compartilhados



Tanenbaum: Figura 6.19

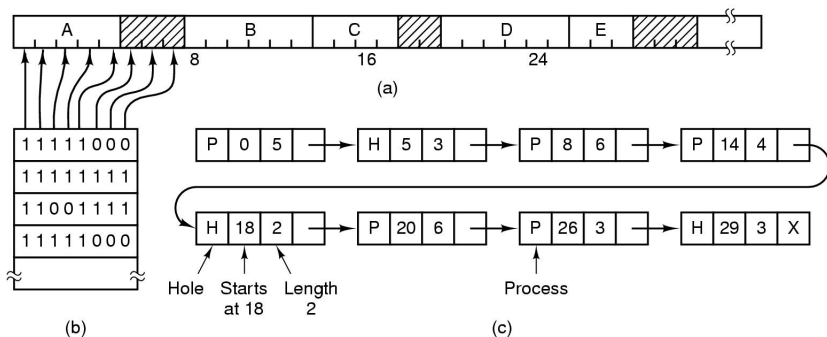
Lista de livres e bitmaps



Tanenbaum: Figura 6.21

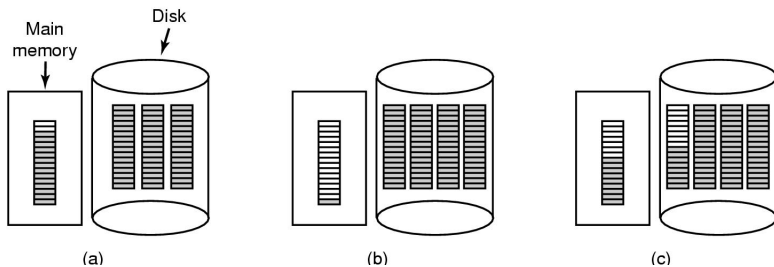
Bitmaps e lista de livres

Relembrando gerência de memória...



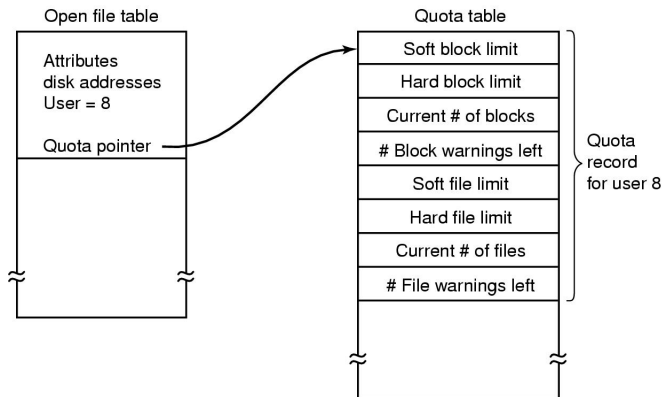
Tanenbaum: Figura 4.7

Lista de livres em memória



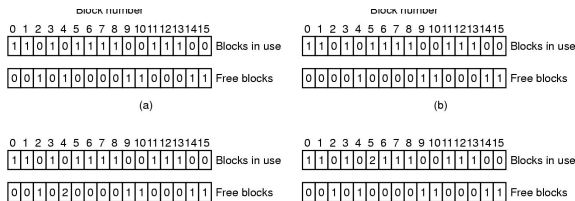
Tanenbaum: Figura 6.22

Gerência de quotas



Tanenbaum: Figura 6.23

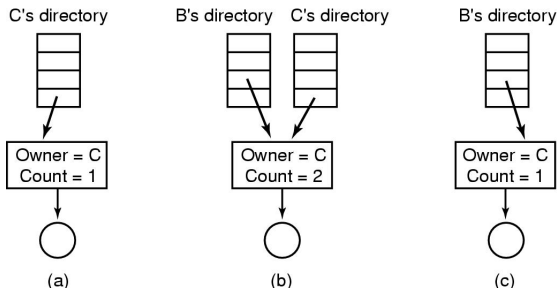
Consistência do sistema de arquivos



Tanenbaum: Figura 6.26

- (a) consistente
- (b) bloco faltando
- (c) duplicação na lista de livres
- (d) duplicação nos dados

Consistência do sistema de arquivos



Tanenbaum: Figura 6.19

Verificar se todos os contadores estão corretos

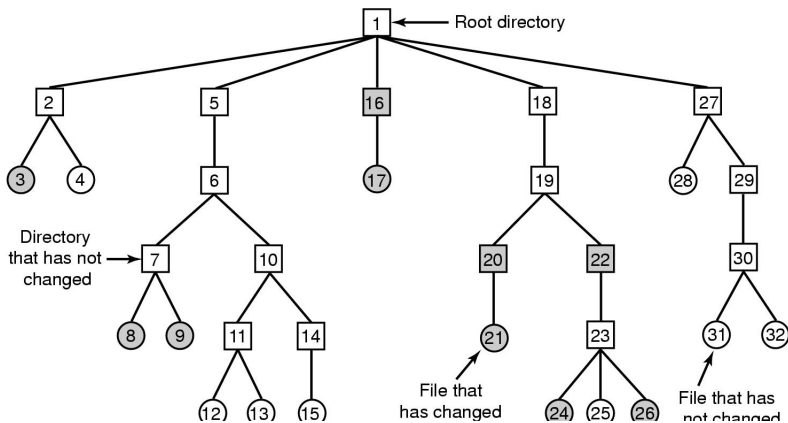
Cópias de segurança

- Dump físico
 - Cópia “total” disco
 - Simples e rápida
 - Blocos livre são copiados?
 - Gerência de blocos defeituosos

Cópias de segurança

- Cópias lógicas
- O que não copiar?
 - Arquivos de instalação do sistema
 - Arquivos /dev
 - Arquivos temporários
- Cópias incrementais

Cópia incremental

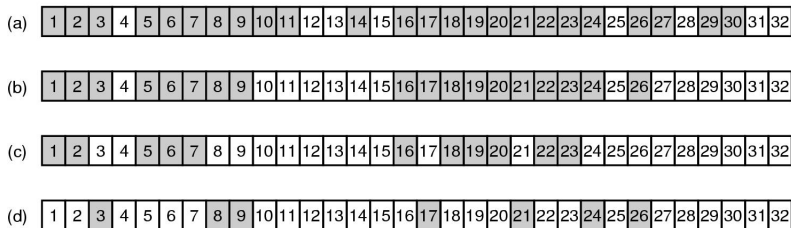


Tanenbaum: Figura 6.24

Cópia incremental

- Mapa de bits representando i-nodes
- Fase 1: marca todos os arquivos modificados e todos os diretórios.
- Fase 2: desmarca todos os diretórios sem arquivos ou sub-diretórios modificados
- Fase 3: varre os i-nodes e copia os diretórios e seus atributos
- Fase 4: copia os arquivos

Mapas de i-nodes e fases

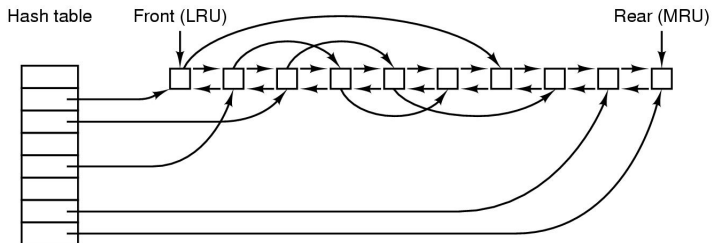


Tanenbaum: Figura 6.25

Restauração de arquivos

- Cria-se um sistema de arquivos vazio
- Cópia completa mais antiga é restaurada
- Cópias incrementais são restauradas
- Complicações
 - Hard links
 - Arquivos com lacunas (e.g., core)

Caching

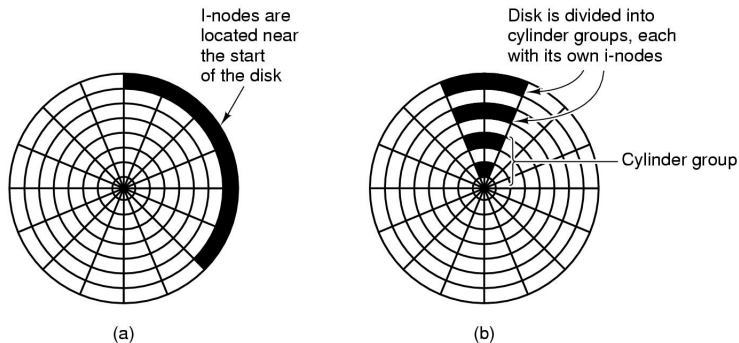


Tanenbaum: Figura 6.27

Block Read Ahead

- Lê um bloco antes de ele ser solicitado
- Acesso seqüencial
- Acesso aleatório

Distribuição da informação no disco



Tanenbaum: Figura 6.28