MC514

Sistemas Operacionais:

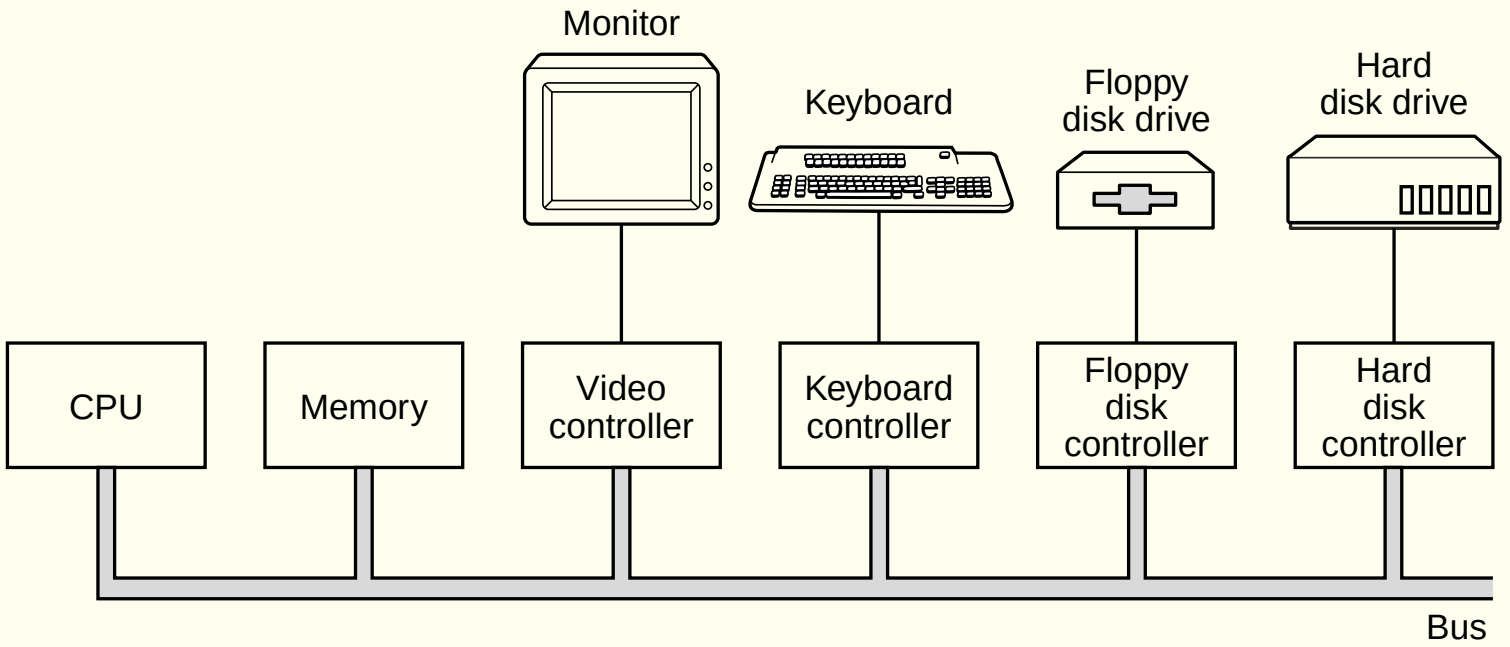
Teoria e Prática

1s2009

Entrada e Saída

Pipes

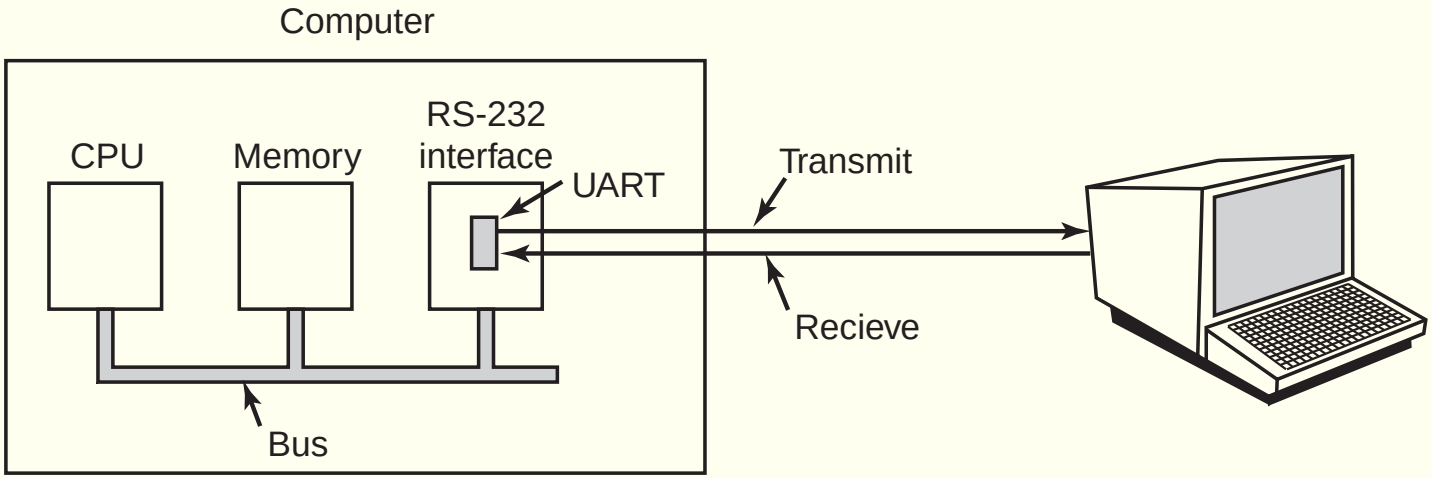
Dispositivos de I/O e controladores



Tanenbaum: Figura 1.5

O sistema operacional deve interagir com os controladores

Character device

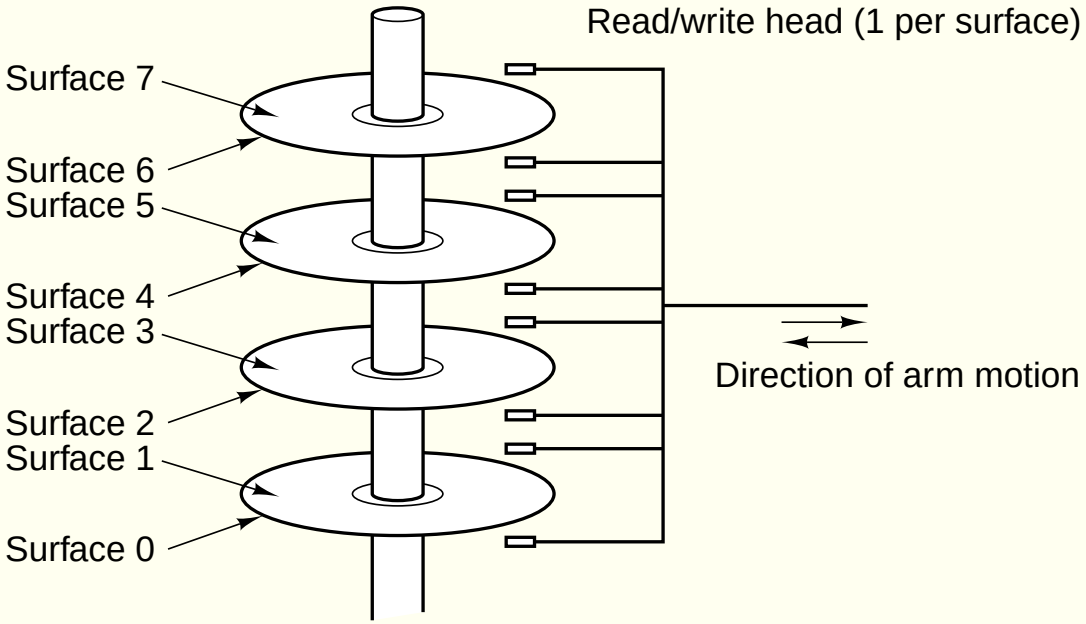


Tanenbaum: Figura 5.34

Acesso sequencial, caracter a caracter

Execute ls -l /dev

Block device

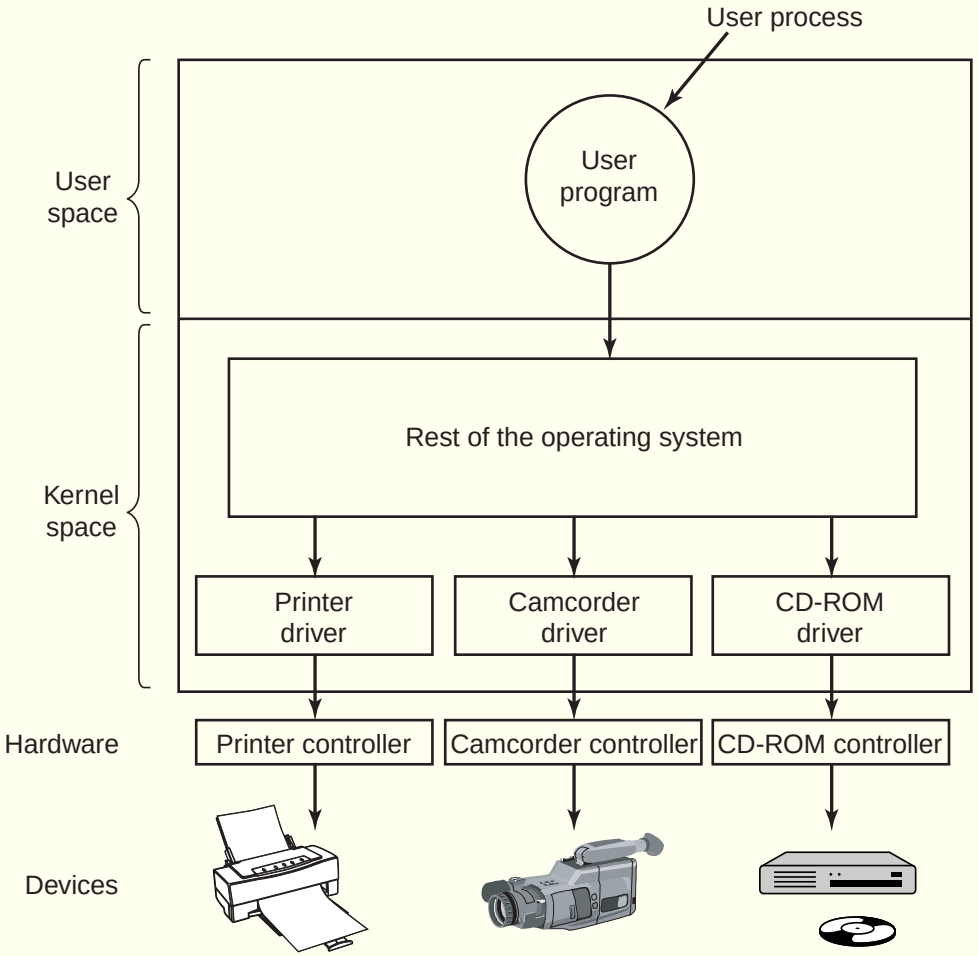


Tanenbaum: Figura 1.8

Acesso não sequencial a blocos de informação

Execute ls -l /dev

Device drivers



Tanenbaum: Figura 5.11

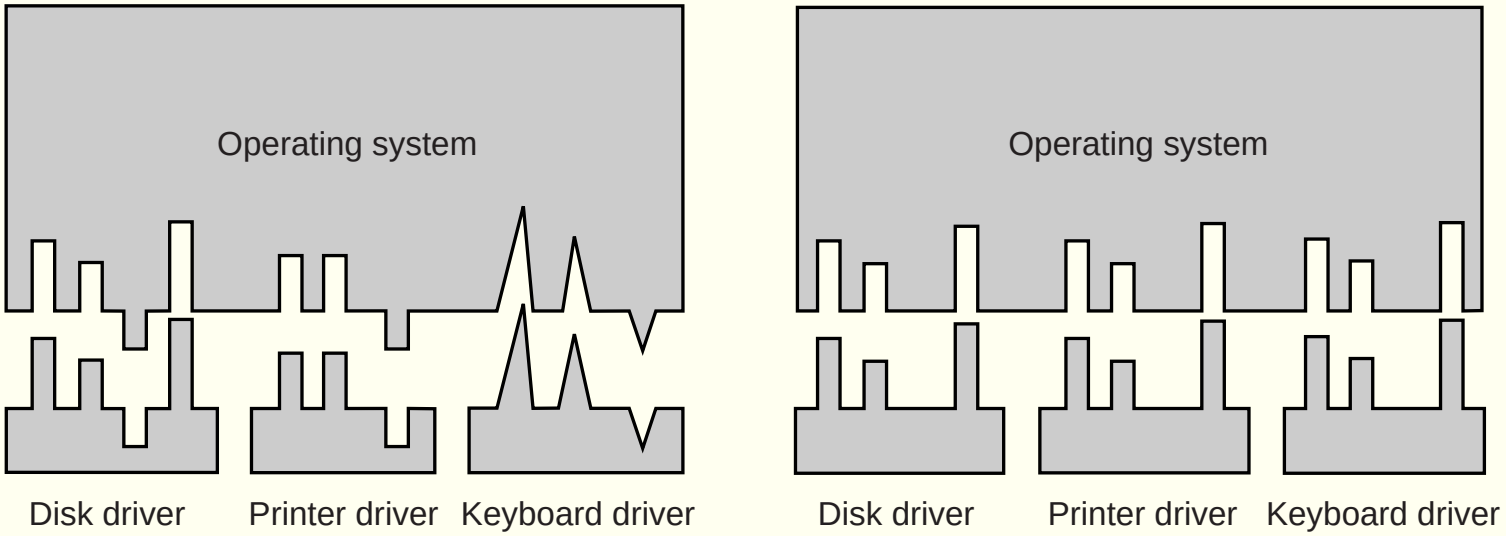
Device drivers

- Software que "conversa" com o controlador
- Os fabricantes devem fornecer device drivers para os sistemas operacionais
- Como acoplar um device driver ao kernel:
  - relink e reboot
  - entrada em um arquivo e reboot
  - on-the-fly

    veja os comandos `lsmod` e `modprobe`

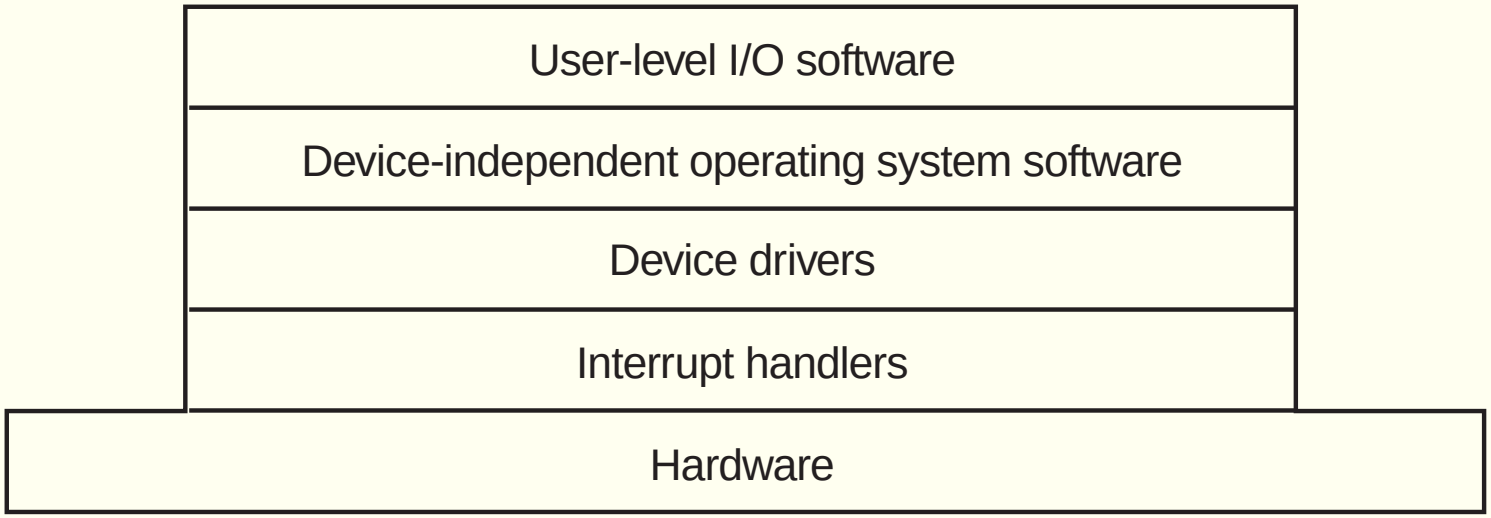
Device drivers

Sem ou com uma interface padrão



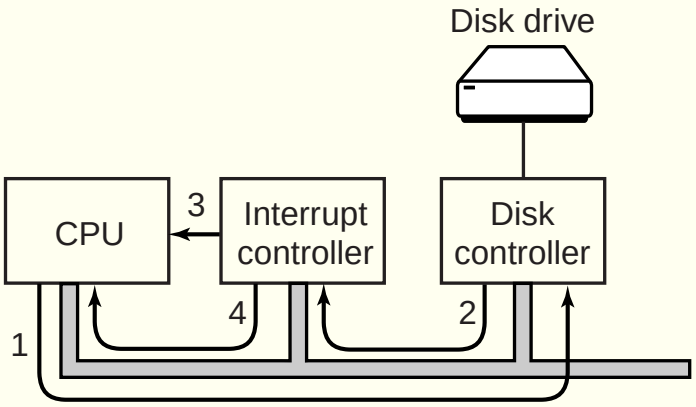
Tanenbaum: Figura 5.13

Camadas de software

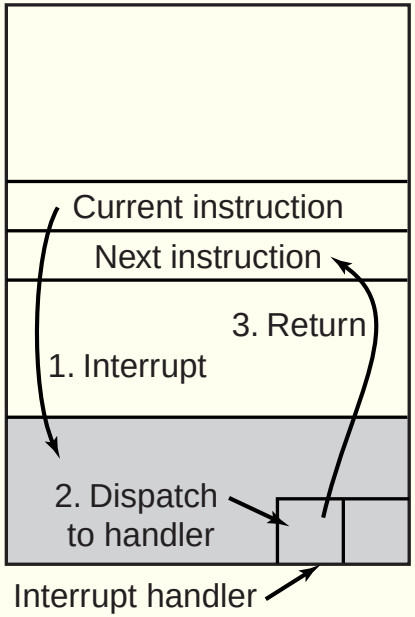


Tanenbaum: Figura 5.10

Tratamento de interrupções



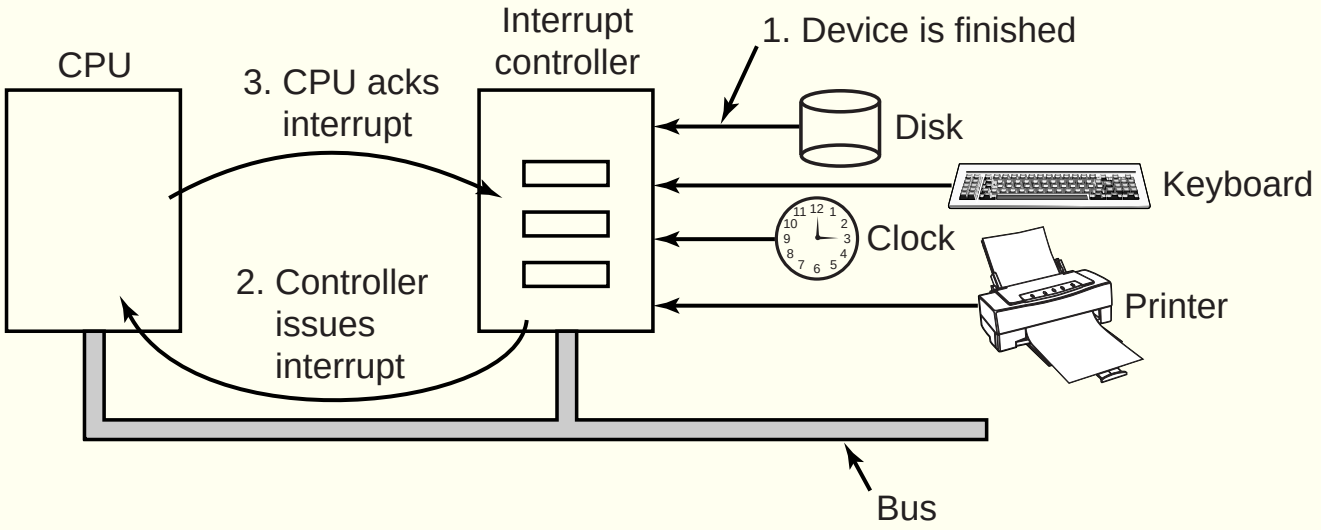
(a)



(b)

Tanenbaum: Figura 1.10

Tratamento de interrupções



Tanenbaum: Figura 5.5

Como programar os dispositivos?

- Instruções especiais

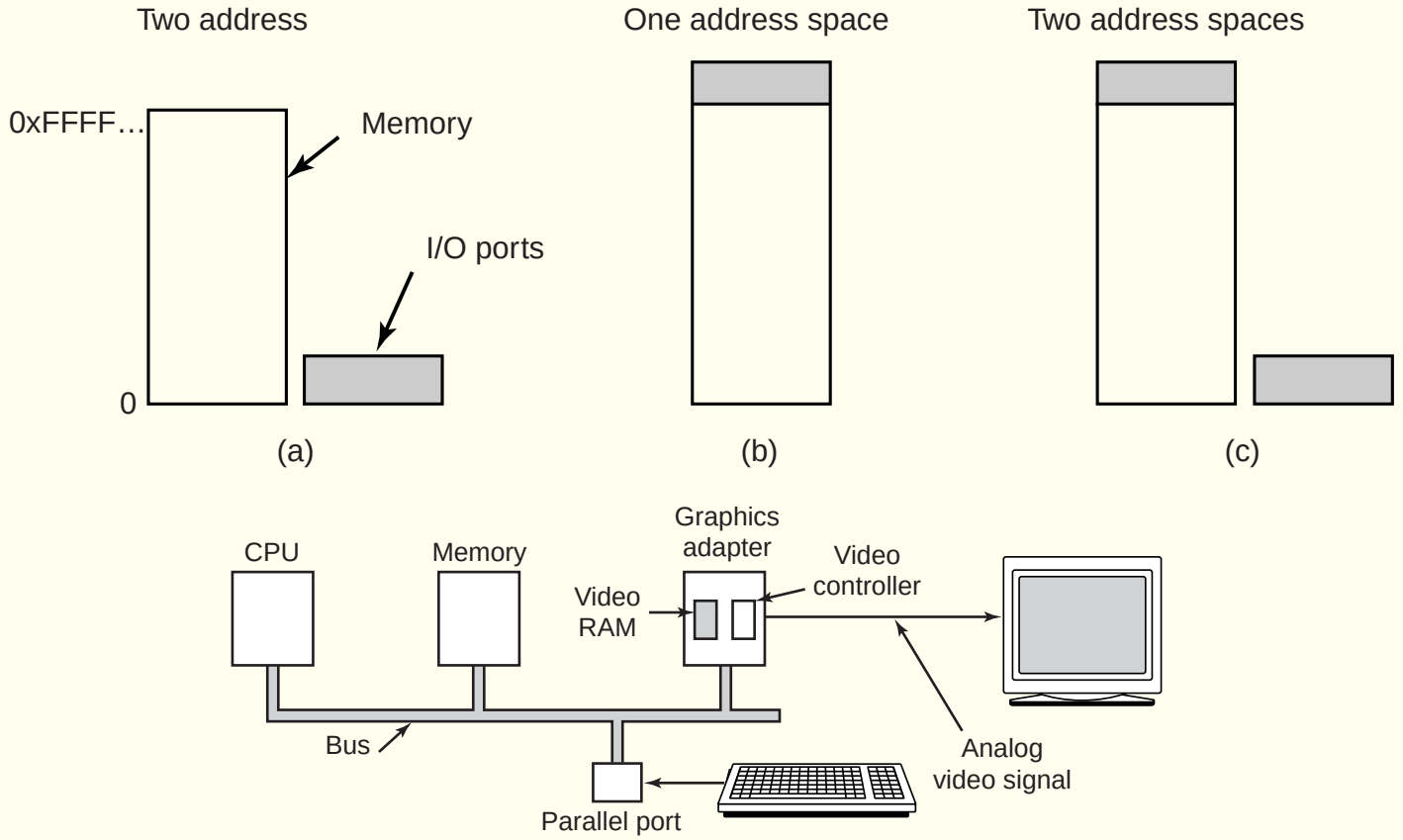
  ```
  IN REG, PORT
  OUT PORT, REG
  ```

- Memory-mapped I/O

  ```
  MOV REG, ADDR
  ```

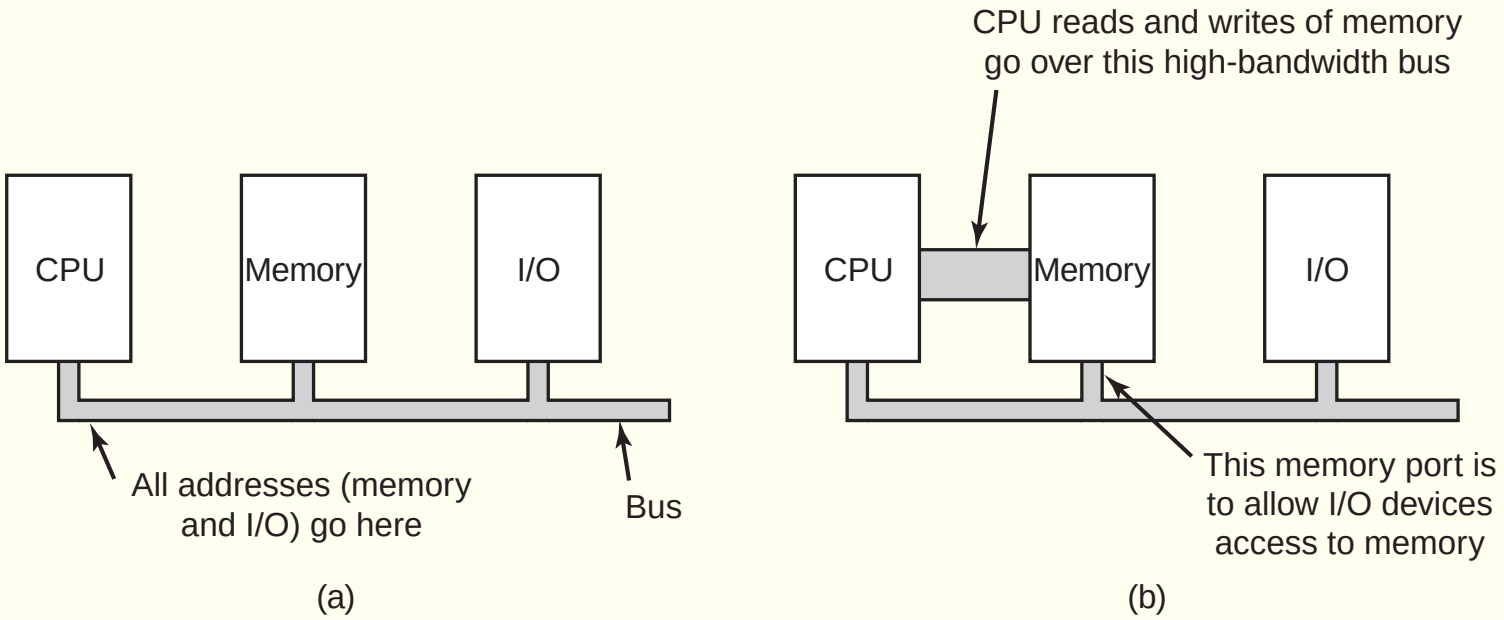
  Conforme o valor de ADDR, a instrução MOV fará acesso a uma palavra de memória ou dispositivo

Como programar os dispositivos?



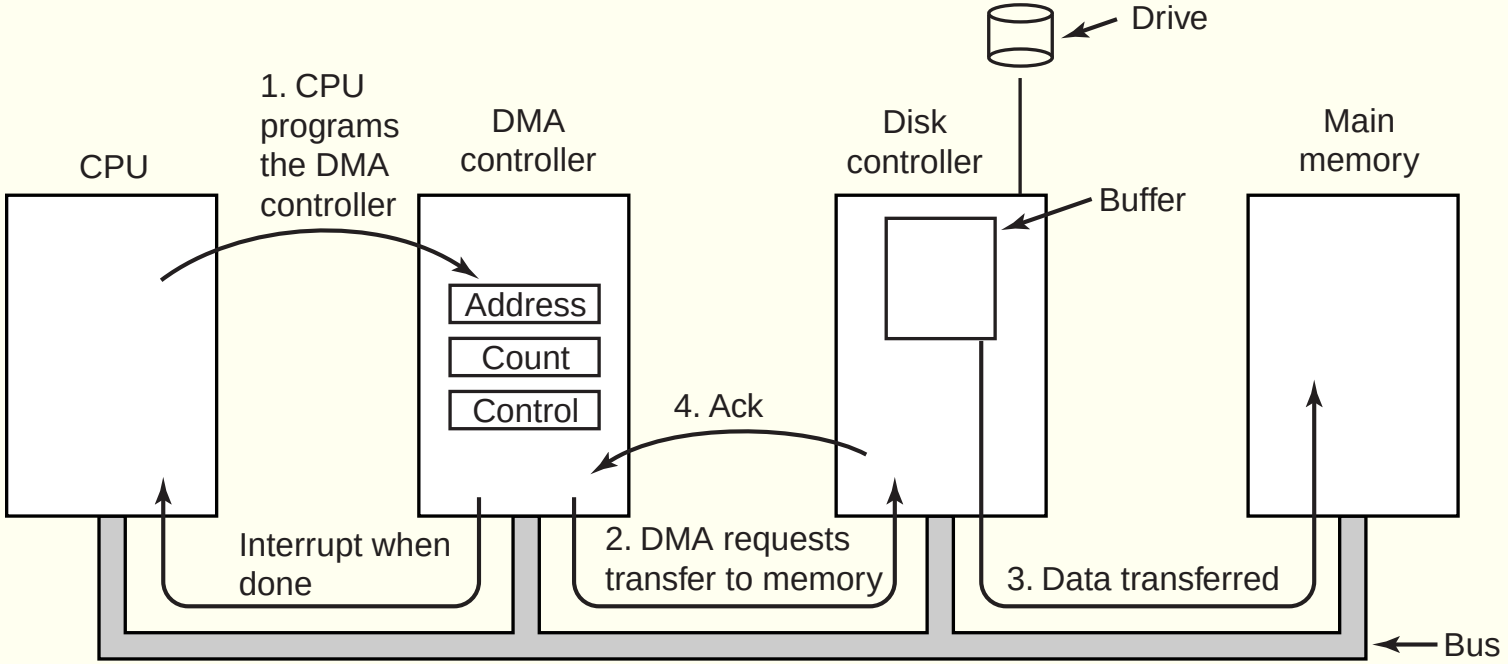
Tanenbaum: Figuras 5.2 e 5-38

Barramento simples e dual



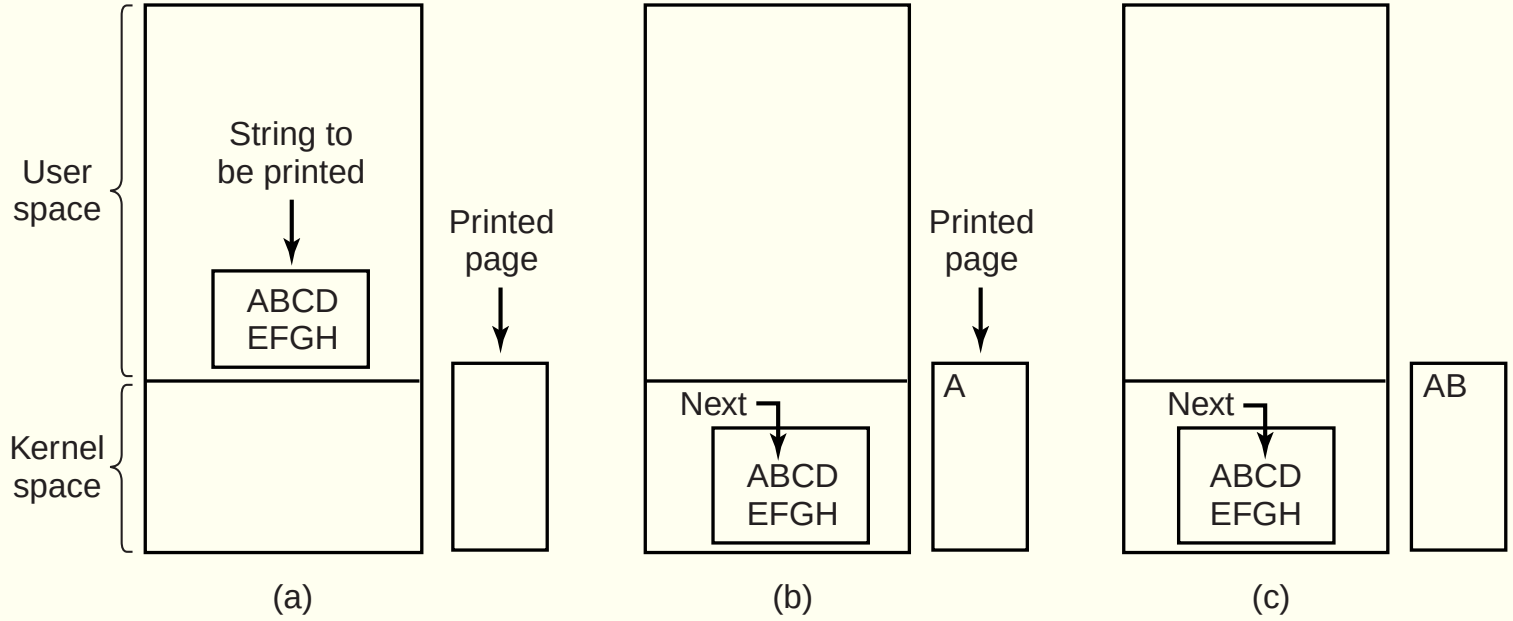
Tanenbaum: Figura 5.3

Direct Memory Access (DMA)



Tanenbaum: Figura 5.4

Imprimindo uma string



Tanenbaum: Figura 5.6

Imprimindo uma string

Programmed I/O

```
copy_from_user(buffer, p, count);                /* p is the kernel bufer */
for (i = 0; i < count; i++) {                    /* loop on every character */
      while (*printer_status_reg != READY) ;     /* loop until ready */
      *printer_data_register = p[i];             /* output one character */
}
return_to_user( );
```

Tanenbaum: Figura 5.7

Trecho de código do kernel

Imprimindo uma string

Interrupt-driven I/O

```
copy_from_user(buffer, p, count);
enable_interrupts( );
while (*printer_status_reg != READY) ;
*printer_data_register = p[0];
scheduler( );
```

(a)

```
if (count == 0) {
    unblock_user( );
} else {
    *printer_data_register = p[i];
    count = count – 1;
     i = i + 1;
}
acknowledge_interrupt( );
return_from_interrupt( );
```

(b)

Tanenbaum: Figura 5.8

(a) Trecho de código do kernel

(b) Tratador da interrupção

Imprimindo uma string

DMA

```
copy_from_user(buffer, p, count);
set_up_DMA_controller( );
scheduler( );
```

(a)

```
acknowledge_interrupt( );
unblock_user( );
return_from_interrupt( );
```

(b)

(a) Trecho de código do kernel

(b) Tratador de interrupção

Pipes

```
$ grep xxx log.txt > log-xxx.txt
$ wc -l log-xxx.txt
$ rm log-xxx.txt

$ grep xxx log.txt | wc -l
```

pipe()

`int pipe (int FILEDES[2])`

> The ‘pipe’ function creates a pipe and puts the file descriptors for the reading and writing ends of the pipe (respectively) into ‘FILEDES[0]’ and ‘FILEDES[1]’.

Veja o código: mypipe.c

Pipe com entrada e saída padrão?

```
int dup2(int oldfd, int newfd);
```

dup2 makes newfd be the copy of oldfd, closing newfd first if necessary. After successful return of dup or dup2, the old and new descriptors may be used interchangeably.

Veja o código: mypipe2.c

Processos conectados de maneira transparente

```
$ cm1 <args1> | cmd2 <args2>
```

- A modificação da entrada e saída padrão deve ser feita antes da chamada a execve().
- Veja o código: minishell.c

popen()

```
FILE *popen(const char *command,
            const char *type);
int pclose(FILE *stream);
```

The popen() function opens a process by creating a pipe, forking, and invoking the shell. Since a pipe is by definition unidirectional, the type argument may specify only reading or writing, not both; the resulting stream is correspondingly read-only or write-only.

Veja o código: mypopen.c e mypopen2.c