

# MC542

## Organização de Computadores Teoria e Prática

2006

Prof. Paulo Cesar Centoducatte

[ducatte@ic.unicamp.br](mailto:ducatte@ic.unicamp.br)

[www.ic.unicamp.br/~ducatte](http://www.ic.unicamp.br/~ducatte)

**MC542**

## **Circuitos Lógicos**

### **Representação de Números e Circuitos Aritméticos**

**“Fundamentals of Digital Logic with VHDL  
Design” - (Capítulo 5)**

# Título do Capítulo Abordado

## Sumário

- Representação Posicional
- Adição de Números sem Sinal
- Números com Sinal (sinalizados)
- Somadores Rápidos
- Multiplicação
- Representações de Números Reais

# Representação Posicional

- Representação de Inteiros sem Sinal

$$D = d_{n-1} d_{n-2} \dots d_1 d_0$$

$$V(D) = d_{n-1} \times 10^{n-1} + d_{n-2} \times 10^{n-2} + \dots + d_1 \times 10^1 + d_0 \times 10^0$$

- Sistema Binário:

$$B = b_{n-1} b_{n-2} \dots b_1 b_0$$

$$V(B) = b_{n-1} \times 2^{n-1} + b_{n-2} \times 2^{n-2} + \dots + b_1 \times 2^1 + b_0 \times 2^0$$

$$= \sum_{i=0}^{n-1} b_i \times 2^i$$

## Representação Posicional Conversão entre Decimal e Binário

$$V(B) = b_{n-1} \times 2^{n-1} + b_{n-2} \times 2^{n-2} + \dots + b_1 \times 2^1 + b_0 \times 2^0$$

$$V(B) = b_{n-1} \times 2^{n-1} + b_{n-2} \times 2^{n-2} + \dots + b_1 \times 2^1 + b_0$$

$$\frac{V(B)}{2} = b_{n-1} \times 2^{n-2} + b_{n-2} \times 2^{n-3} + \dots + b_1 + \frac{b_0}{2}$$

Conversão de Decimal para Binário: Divisão sucessiva por 2

# Exemplo

Convert  $(857)_{10}$

				Remainder	
$857 \div 2$	$=$	428	1		LSB
$428 \div 2$	$=$	214	0		
$214 \div 2$	$=$	107	0		
$107 \div 2$	$=$	53	1		
$53 \div 2$	$=$	26	1		
$26 \div 2$	$=$	13	0		
$13 \div 2$	$=$	6	1		
$6 \div 2$	$=$	3	0		
$3 \div 2$	$=$	1	1		
$1 \div 2$	$=$	0	1		MSB

Result is  $(1101011001)_2$

# Octal e Hexadecimal

Decimal	Binary	Octal	Hexadecimal
00	00000	00	00
01	00001	01	01
02	00010	02	02
03	00011	03	03
04	00100	04	04
05	00101	05	05
06	00110	06	06
07	00111	07	07
08	01000	10	08
09	01001	11	09
10	01010	12	0A
11	01011	13	0B
12	01100	14	0C
13	01101	15	0D
14	01110	16	0E
15	01111	17	0F
16	10000	20	10
17	10001	21	11
18	10010	22	12

# Adição de Números sem Sinal

$$\begin{array}{r} 0\ 1\ 0\ 1 \\ +\ 0\ 0\ 1\ 1 \\ \hline 0\ 1\ 1\ 0 \end{array}$$

$$\begin{array}{r} 0\ x \\ +\ 0\ y \\ \hline c\ s \end{array}$$

		Carry	Soma
x	y	c	s
0	0	0	0
0	1	0	1
1	0	0	1
1	1	1	0

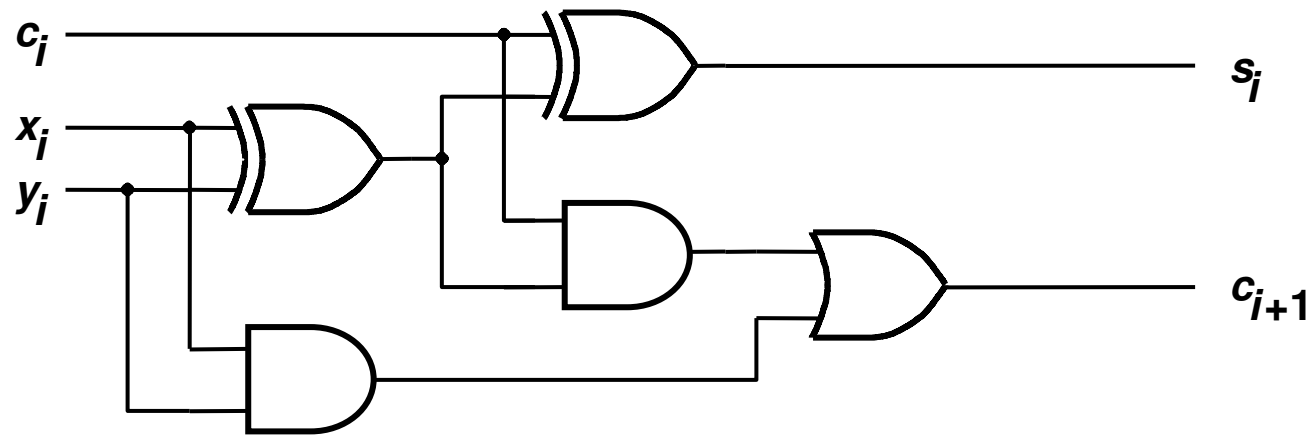
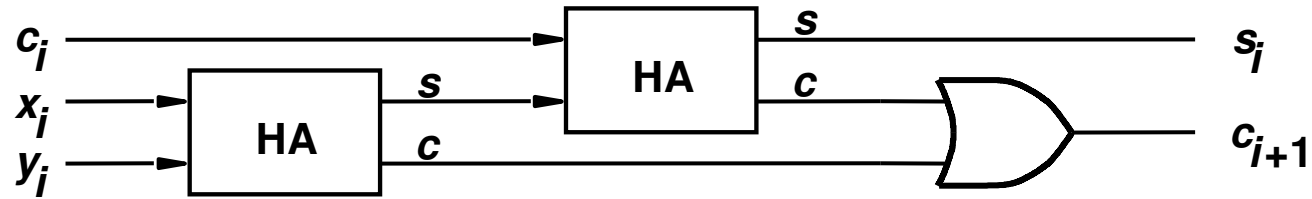


# Adição de Números sem Sinal

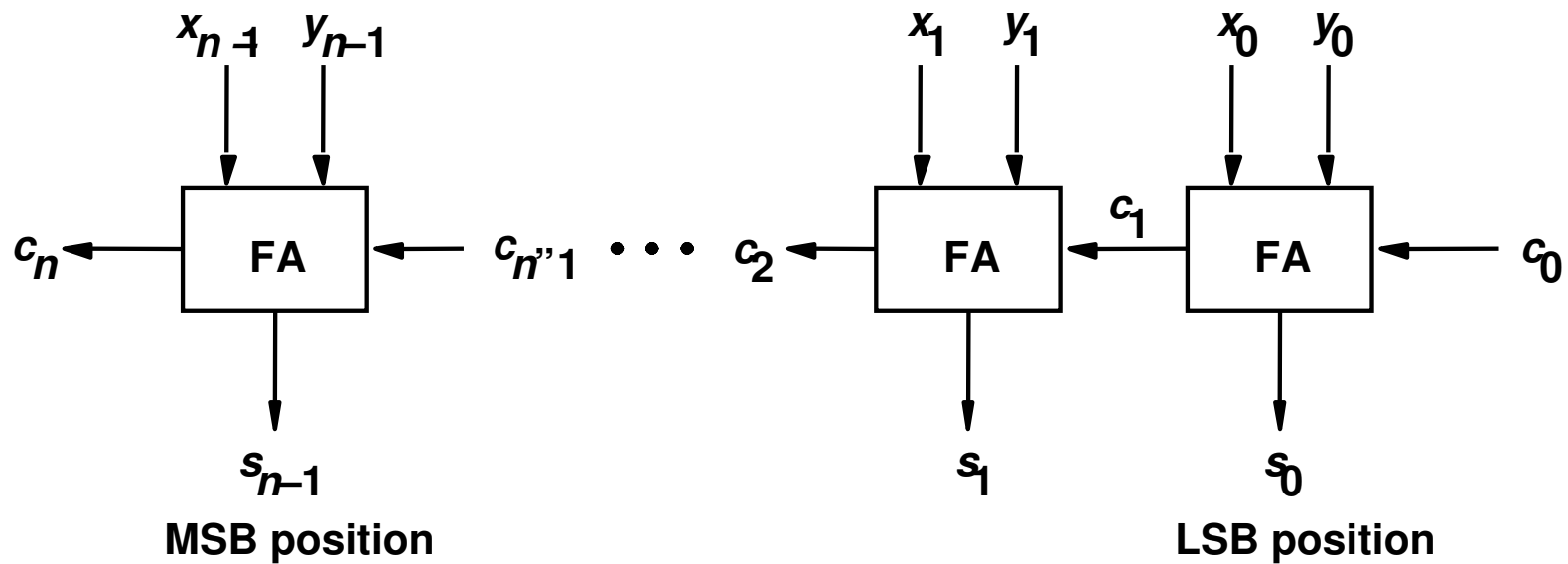
$$\begin{array}{r} X = x_4 x_3 x_2 x_1 x_0 \quad 01111 \quad (15)_{10} \\ + Y = y_4 y_3 y_2 y_1 y_0 \quad 01010 \quad (10)_{10} \\ \hline \quad 1110 \quad \leftarrow \text{Generated carries} \\ \hline S = s_4 s_3 s_2 s_1 s_0 \quad 11001 \quad (25)_{10} \end{array}$$

# Adição de Números sem Sinal

## Somador Completo a partir de $\frac{1}{2}$ Somador

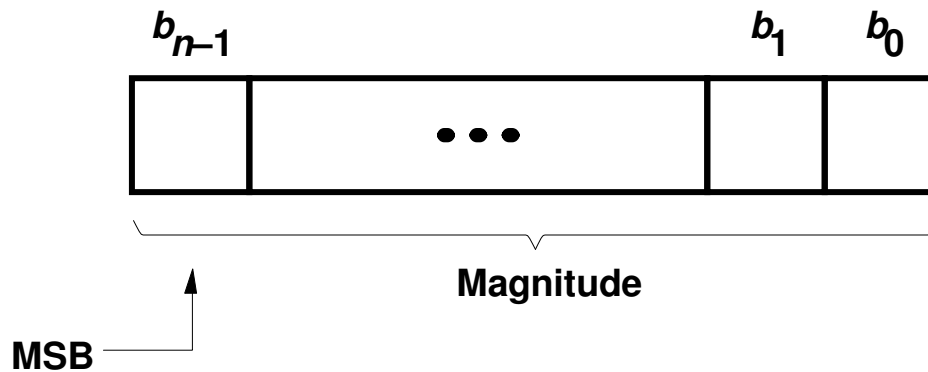


# Somador Ripple-Carry

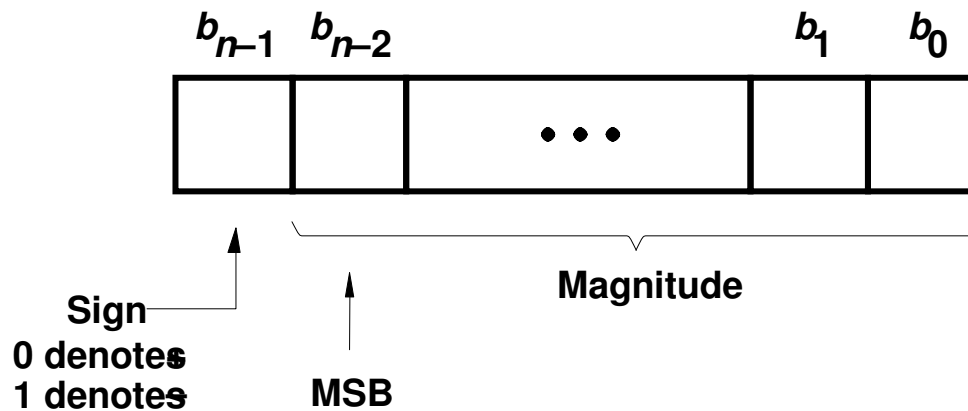


# Representação de Números Negativos

## Sinal e Magnitude



Número sem Sinal



Número com Sinal

# Representação de Números Negativos

## Complemento de 1

Em complemento de "Um" o número negativo  $K$ , com  $n$ -bits, é obtido subtraindo seu positivo  $P$  de  $2^n - 1$

$$K = (2^n - 1) - P$$

Exemplo: se  $n = 4$   
então:

$$\begin{aligned} K &= (2^4 - 1) - P \\ K &= (16 - 1) - P \\ K &= (1111)_2 - P \end{aligned}$$

$$\begin{aligned} K = -7 &\quad \rightarrow \quad P = 7 \\ 7 &= (0111)_2 \\ -7 &= (1111)_2 - (0111)_2 \\ -7 &= (1000)_2 \end{aligned}$$

# Representação de Números Negativos

## Complemento de 2

Em complemento de "Dois" o número negativo  $K$ , com  $n$ -bits, é obtido subtraindo seu positivo  $P$  de  $2^n$

$$K = 2^n - P \quad \longrightarrow$$

$$K = (2^n - 1) + 1 - P$$

$$K = (2^n - 1) - P + 1$$

Exemplo: se  $n = 4$   
então:

$$K = 2^4 - P$$

$$K = 16 - P$$

$$K = (10000)_2 - P$$

$$K = -7 \quad \rightarrow \quad P = 7$$

$$7 = (0111)_2$$

$$-7 = (10000)_2 - (0111)_2$$

$$-7 = (1001)_2$$

# Representação de Números Negativos

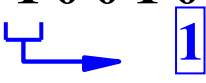
$b_3b_2b_1b_0$	Sign and magnitude	1's complement	2's complement
0111	+7	+7	+7
0110	+6	+6	+6
0101	+5	+5	+5
0100	+4	+4	+4
0011	+3	+3	+3
0010	+2	+2	+2
0001	+1	+1	+1
0000	+0	+0	+0
1000	-0	-7	-8
1001	-1	-6	-7
1010	-2	-5	-6
1011	-3	-4	-5
1100	-4	-3	-4
1101	-5	-2	-3
1110	-6	-1	-2
1111	-7	-0	-1

# Adição e Subtração Complemento de 1

$$\begin{array}{r}
 (+5) \quad 0101 \\
 +(+2) \quad +0010 \\
 \hline
 (+7) \quad 0111
 \end{array}$$

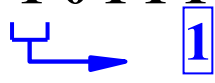
$$\begin{array}{r}
 (-5) \quad 1010 \\
 +(+2) \quad +0010 \\
 \hline
 (-3) \quad 1100
 \end{array}$$

$$\begin{array}{r}
 (+5) \quad 0101 \\
 +(-2) \quad +1101 \\
 \hline
 (+3) \quad 10010
 \end{array}$$



$$\begin{array}{r}
 \boxed{0011}
 \end{array}$$

$$\begin{array}{r}
 (-5) \quad 1010 \\
 +(-2) \quad +1101 \\
 \hline
 (-7) \quad 10111
 \end{array}$$



$$\begin{array}{r}
 \boxed{1000}
 \end{array}$$



# Adição e Subtração Complemento de 2

$$\begin{array}{r}
 (+5) \quad 0101 \\
 + (+2) \quad 0010 \\
 \hline
 (+7) \quad 0111
 \end{array}$$

$$\begin{array}{r}
 (-5) \quad 1011 \\
 + (+2) \quad 0010 \\
 \hline
 (-3) \quad 1101
 \end{array}$$

$$\begin{array}{r}
 (+5) \quad 0101 \\
 + (-2) \quad 1110 \\
 \hline
 (+3) \quad 10011
 \end{array}$$

↑  
ignore

$$\begin{array}{r}
 (-5) \quad 1011 \\
 + (-2) \quad 1110 \\
 \hline
 (-7) \quad 11001
 \end{array}$$

↑  
ignore

# Subtração em Complemento de 2

$$\begin{array}{r} (+5) \quad 0101 \\ - (+2) \quad \underline{0010} \\ \hline (+3) \end{array} \quad \Rightarrow \quad \begin{array}{r} 0101 \\ + 1110 \\ \hline 10011 \end{array}$$

↑  
ignore

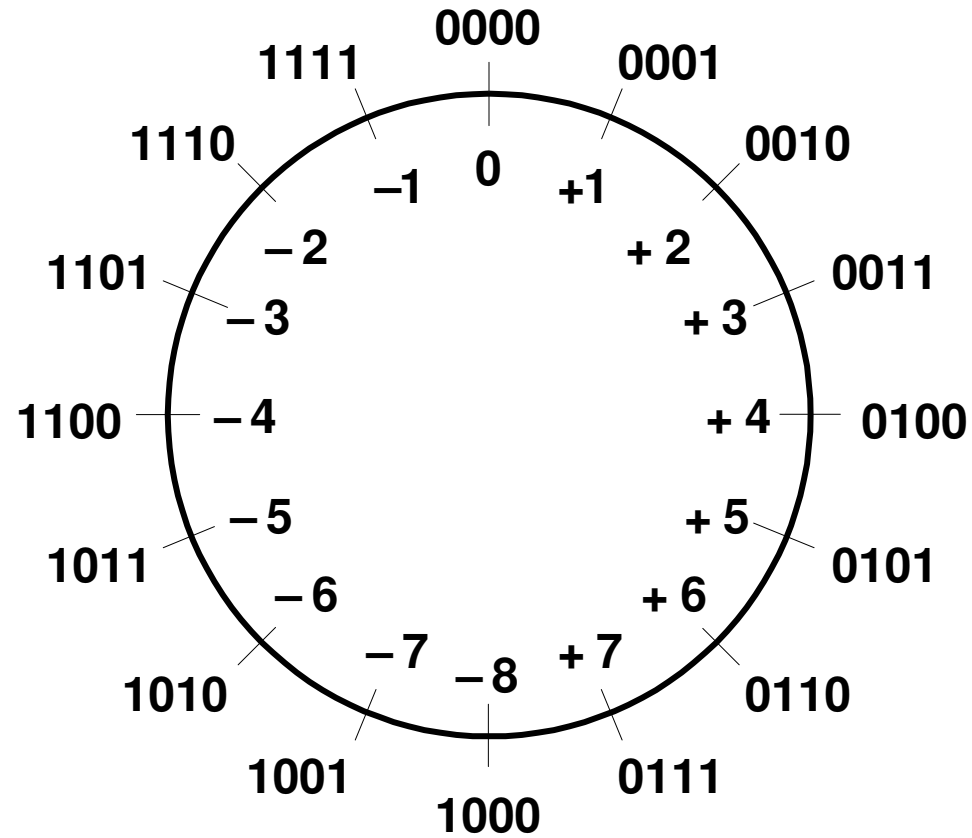
$$\begin{array}{r} (-5) \quad 1011 \\ - (+2) \quad \underline{0010} \\ \hline (-7) \end{array} \quad \Rightarrow \quad \begin{array}{r} 1011 \\ + 1110 \\ \hline 11001 \end{array}$$

↑  
ignore

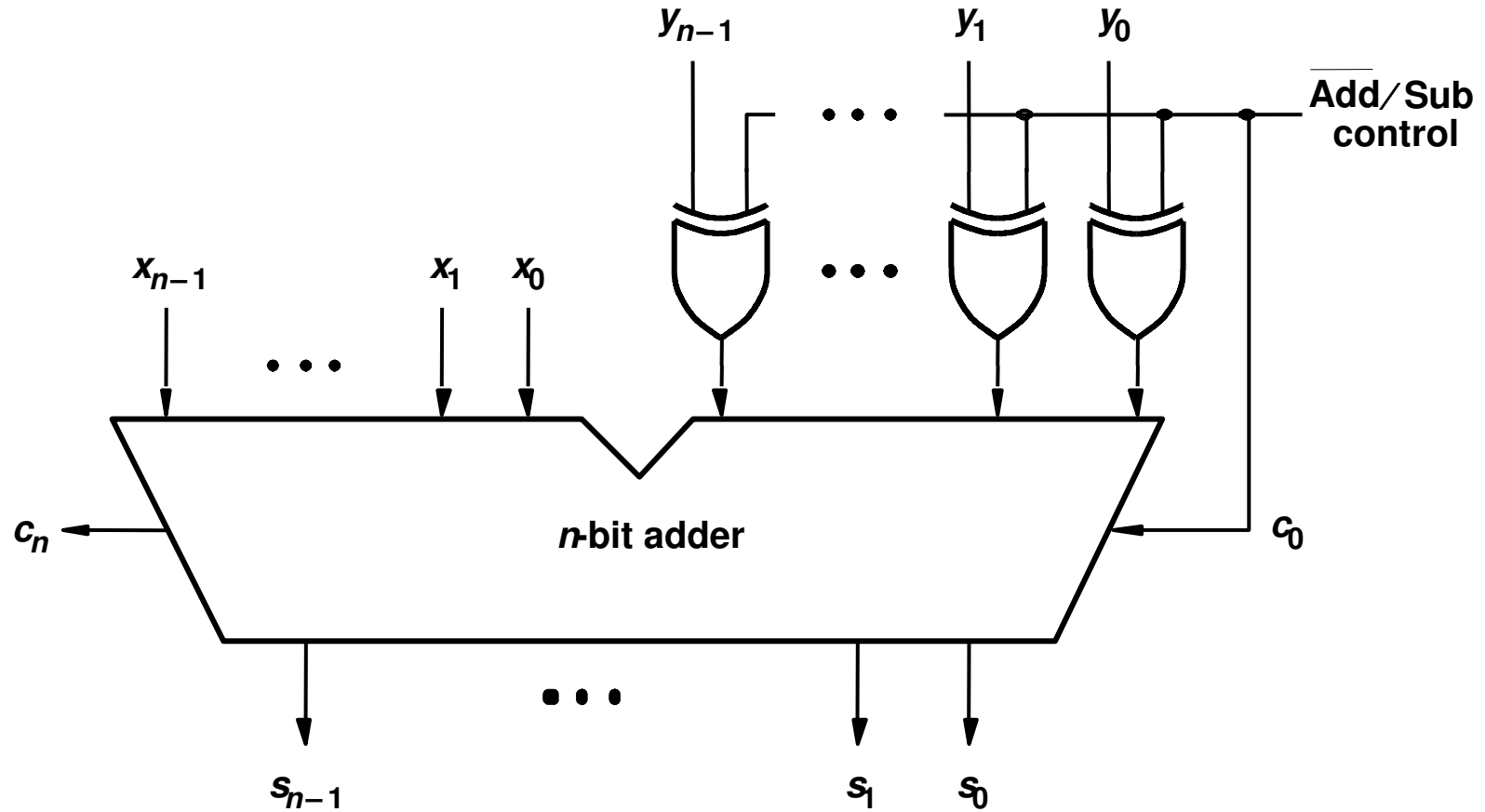
$$\begin{array}{r} (+5) \quad 0101 \\ - (-2) \quad \underline{1110} \\ \hline (+7) \end{array} \quad \Rightarrow \quad \begin{array}{r} 0101 \\ + 0010 \\ \hline 0111 \end{array}$$

$$\begin{array}{r} (-5) \quad 1011 \\ - (-2) \quad \underline{1110} \\ \hline (-3) \end{array} \quad \Rightarrow \quad \begin{array}{r} 1011 \\ + 0010 \\ \hline 1101 \end{array}$$

# Complemento de 2



# Unidade Somador-Subtrator



# Overflow

Quando há overflow?

Como detectar se houve overflow?

$$\begin{array}{r} (+7) \quad 0111 \\ + (+2) \quad +0010 \\ \hline (+9) \quad 1001 \\ c_4 = 0 \\ c_3 = 1 \end{array}$$

$$\begin{array}{r} (-7) \quad 1001 \\ + (+2) \quad +0010 \\ \hline (-5) \quad 1011 \\ c_4 = 0 \\ c_3 = 0 \end{array}$$

$$\begin{array}{r} (+7) \quad 0111 \\ + (-2) \quad +1110 \\ \hline (+5) \quad 10101 \\ c_4 = 1 \\ c_3 = 1 \end{array}$$

$$\begin{array}{r} (-7) \quad 1001 \\ + (-2) \quad +1110 \\ \hline (-9) \quad 10111 \\ c_4 = 1 \\ c_3 = 0 \end{array}$$

# Somadores Rápidos

## Carry-Lookahead

- Desempenho do somador Ripple Carry de n-bits
  - Se cada somador completo tem um atraso de  $\Delta t$
  - Então, o atraso de todo o somador será de  $n \Delta t$

## Somadores Rápidos Carry-Lookahead

$$C_{i+1} = x_i y_i + x_i c_i + y_i c_i$$

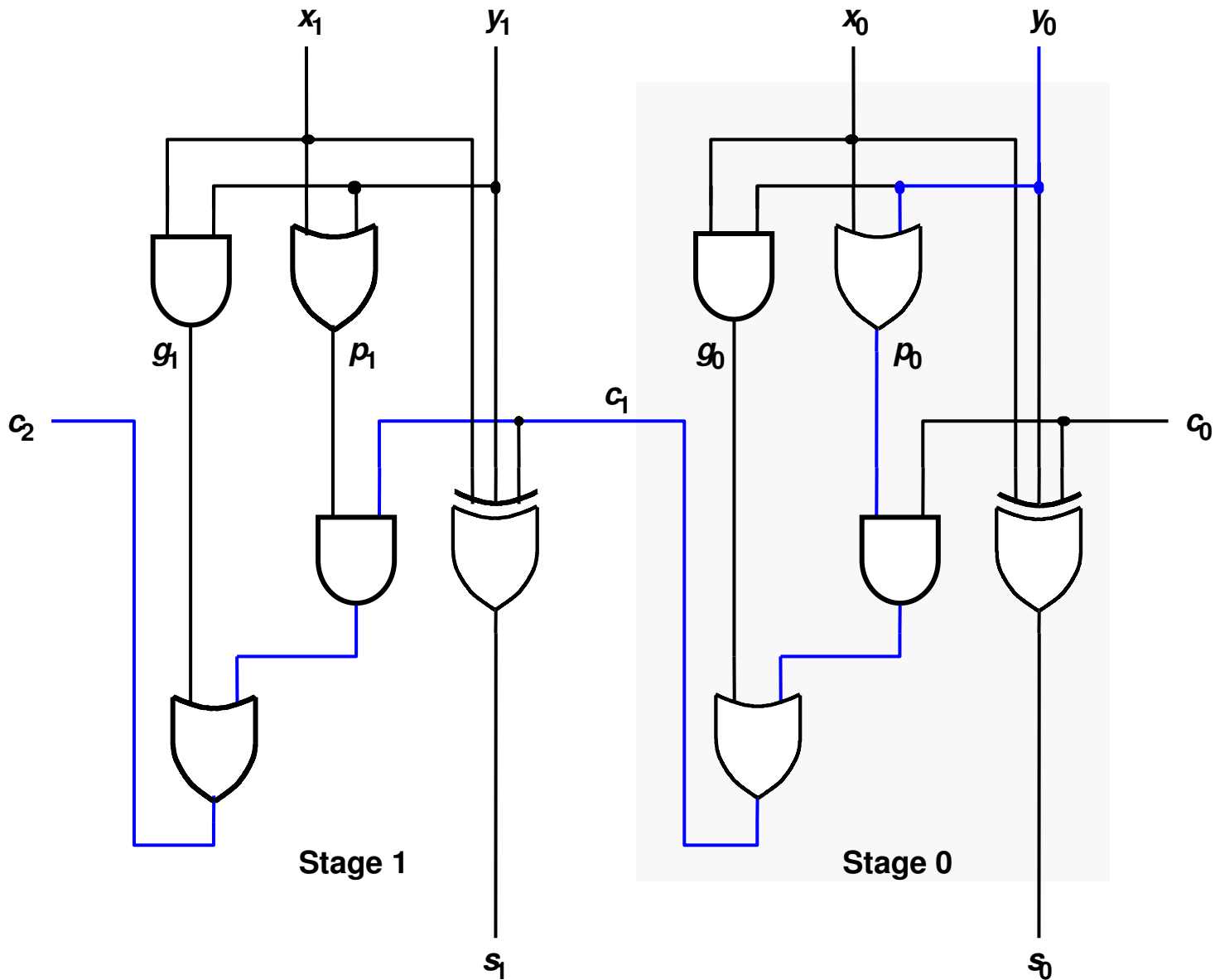
$$C_{i+1} = x_i y_i + (x_i + y_i) c_i \quad \Rightarrow \quad C_{i+1} = g_i + p_i c_i$$

$$C_{i+1} = g_i + p_i (g_{i-1} + p_{i-1} c_{i-1})$$

$$C_{i+1} = g_i + p_i g_{i-1} + p_i p_{i-1} c_{i-1}$$

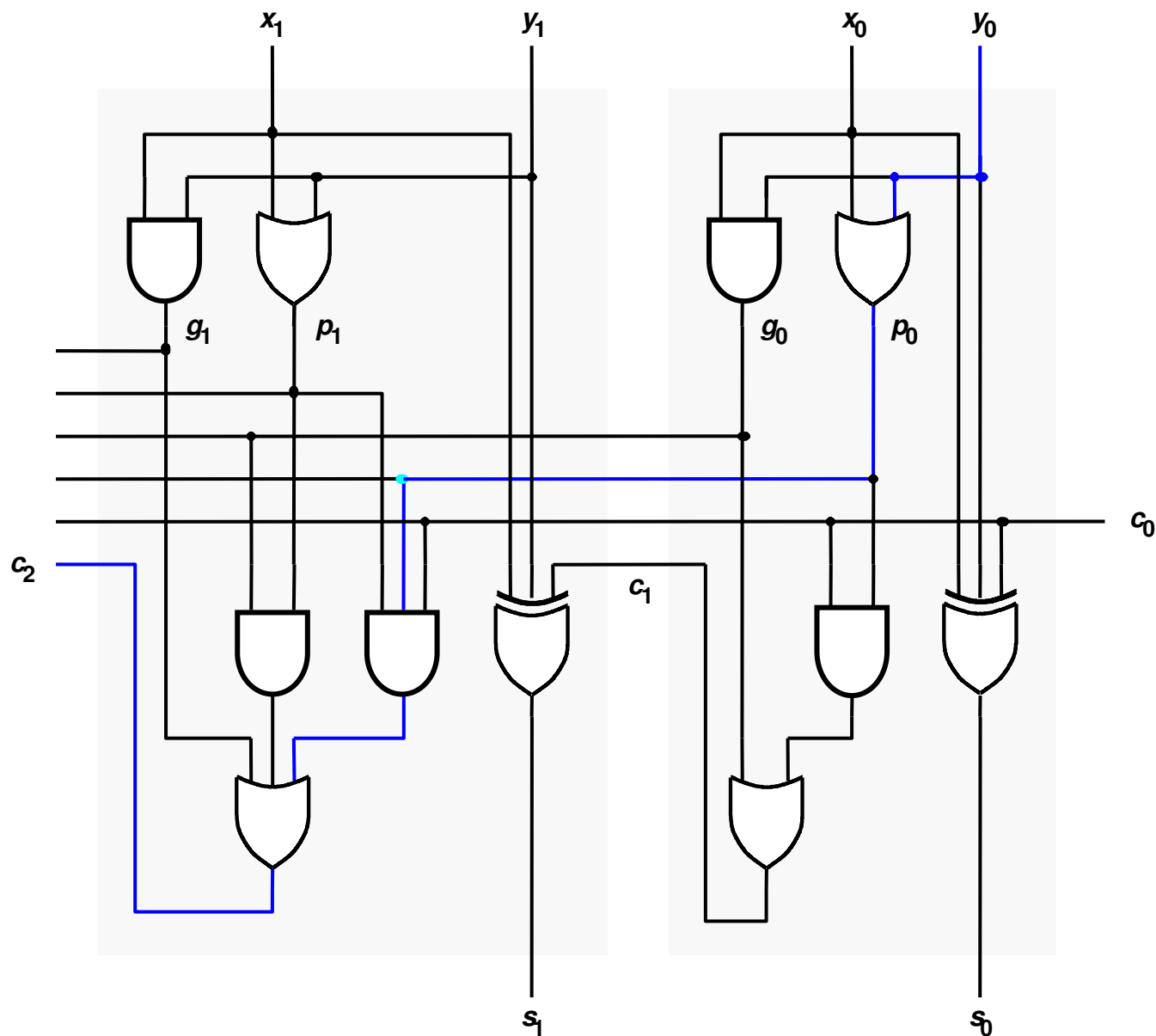
# Somadores Rápidos

## Ripple Carry em termos de g e p

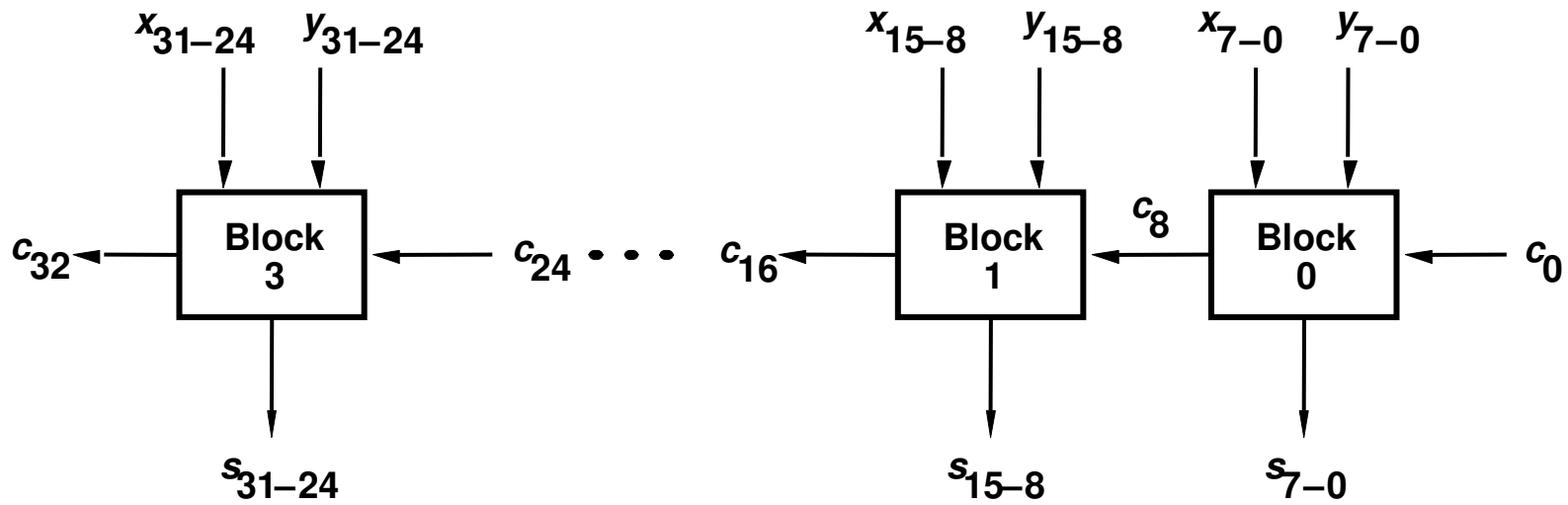




# Somadores Rápidos Carry-Lookahead



# Somadores Rápidos Carry-Lookahead





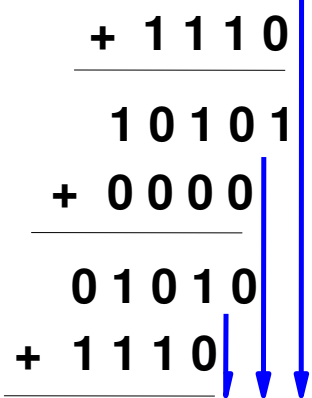
# Multiplicação

<b>Multiplicand M</b>	<b>(14)</b>	<b>1 1 1 0</b>
<b>Multiplier Q</b>	<b>(11)</b>	<b>× 1 0 1 1</b>
		<hr/>
		<b>1 1 1 0</b>
		<b>1 1 1 0</b>
		<b>0 0 0 0</b>
		<b>1 1 1 0</b>
		<hr/>
<b>Product P</b>	<b>(154)</b>	<b>1 0 0 1 1 0 1 0</b>

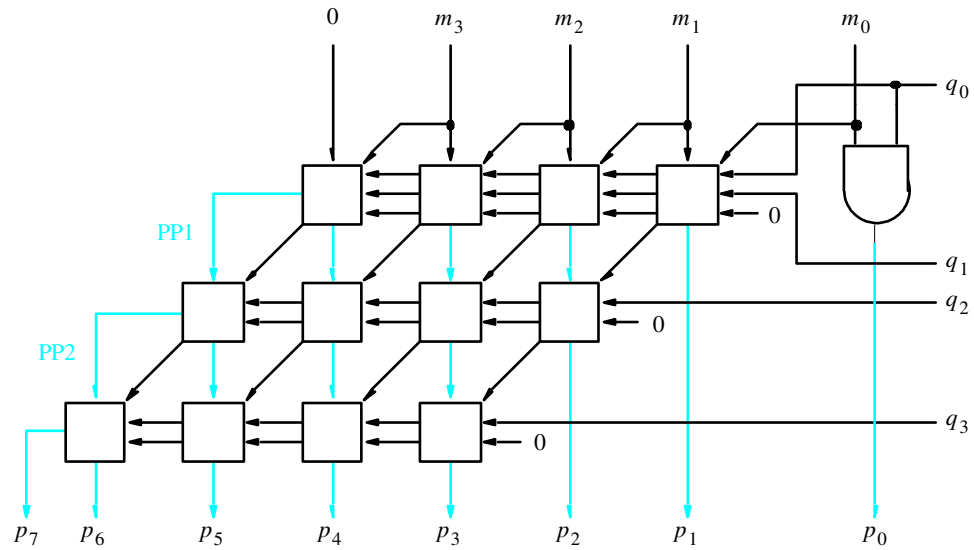
**Multiplicação: lapis e papel**

# Multiplicação

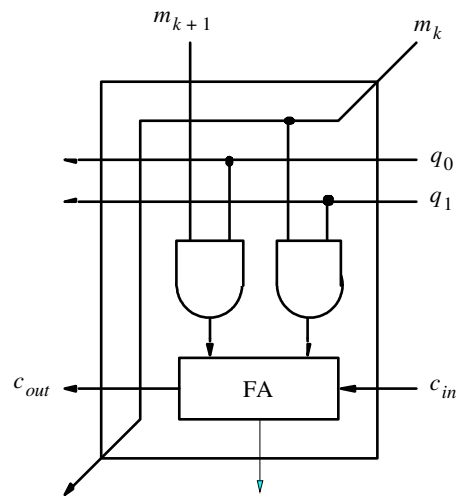
Multiplicand M	(11)	1 1 1 0
Multiplier Q	(14)	× 1 0 1 1
		-----
Partial product 0		1 1 1 0
		+ 1 1 1 0
		-----
Partial product 1		1 0 1 0 1
		+ 0 0 0 0
		-----
Partial product 2		0 1 0 1 0
		+ 1 1 1 0
		-----
Product P	(154)	1 0 0 1 1 0 1 0



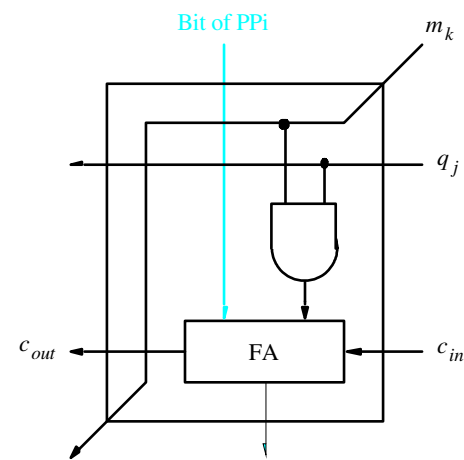
# Multiplicação



(a) Structure of the circuit



(b) A block in the top row



(c) A block in the bottom two rows

# Multiplicação Números com Sinal

Multiplicand M	(+14)	0 1 1 1 0
Multiplier Q	(+11)	x 0 1 0 1 1
Partial product 0		0 0 0 1 1 1 0
		+ 0 0 1 1 1 0
Partial product 1		0 0 1 0 1 0 1
		+ 0 0 0 0 0 0
Partial product 2		0 0 0 1 0 1 0
		+ 0 0 1 1 1 0
Partial product 3		0 0 1 0 0 1 1
		+ 0 0 0 0 0 0
Product P	(+154)	0 0 1 0 0 1 1 0 1 0

(a) Positive multiplicand

# Multiplicação Números com Sinal

Multiplicand M	(-14)	1 0 0 1 0
Multiplier Q	(+11)	× 0 1 0 1 1
Partial product 0		1 1 1 0 0 1 0
		+ 1 1 0 0 1 0
Partial product 1		1 1 0 1 0 1 1
		+ 0 0 0 0 0 0
Partial product 2		1 1 1 0 1 0 1
		+ 1 1 0 0 1 0
Partial product 3		1 1 0 1 1 0 0
		+ 0 0 0 0 0 0
Product P	(-154)	1 1 0 1 1 0 0 1 1 0

(b) Negative multiplicand



# Representações de Números Reais

- Notação Científica:

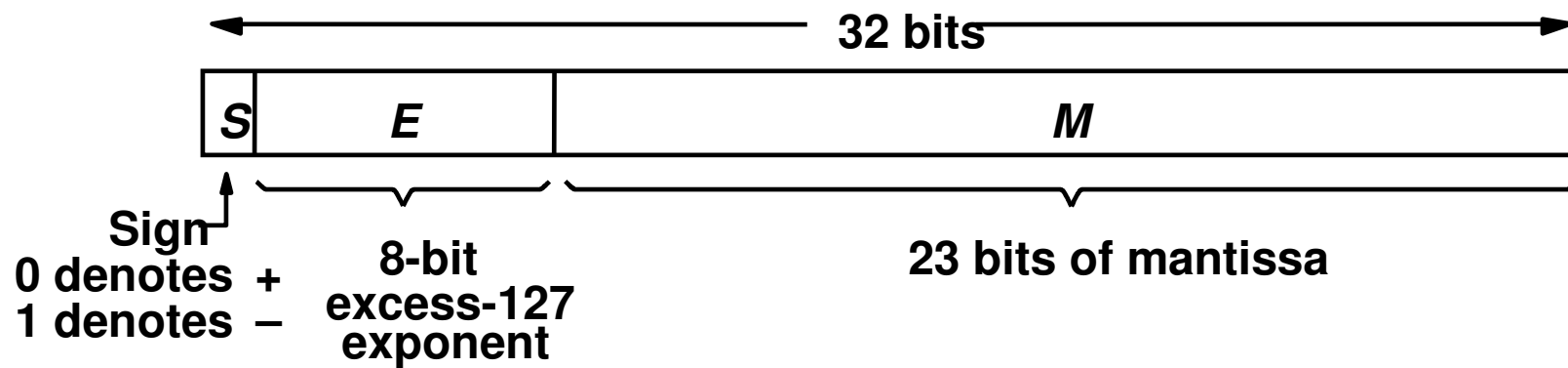
$$\pm \text{Mantissa} \times 10^E$$

$$\text{Mantissa} = x.yyyyyyy$$

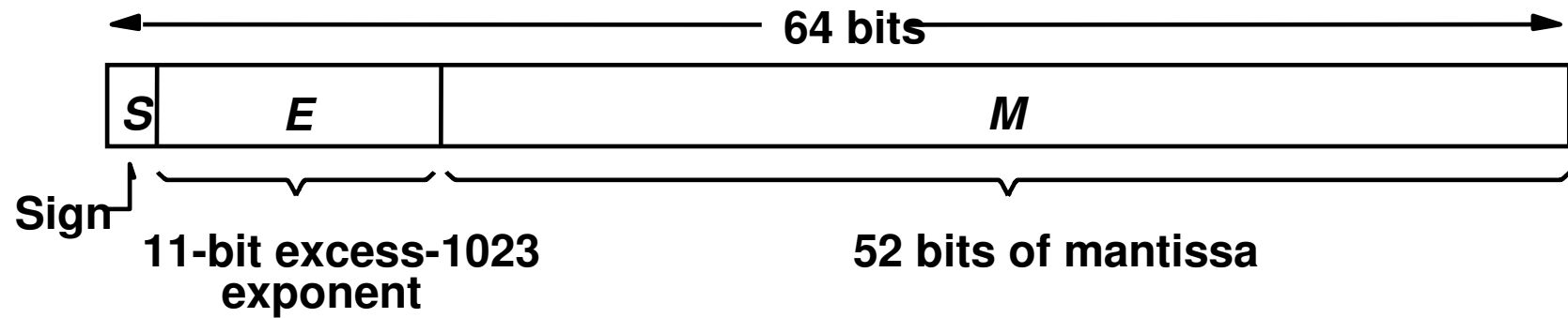
- Se Mantissa não possui zeros a esquerda do ponto decimal -> normalizado
- Padrão IEEE-754
  - Normalizado
  - Bit Escondido

$$(-1)^s \times (1 + \text{Fração}) \times 2^E$$

# IEEE-754 de Precisão Simples



# IEEE-754 de Precisão Dupla



# IEEE-754

## Valores Representados

Exponent	Fraction	Represents
$e = 0 = E_{\text{mim}} - 1$	$f = 0$	$\pm 0$
$e = 0 = E_{\text{mim}} - 1$	$f \neq 0$	$0.f \times 2^{E_{\text{mim}}}$
$E_{\text{mim}} \leq e \leq E_{\text{max}}$		$1.f \times 2^e$
$e = E_{\text{max}} + 1$	$f = 0$	$\pm \infty$
$e = E_{\text{max}} + 1$	$f \neq 0$	NaN





# BCD

## Binary-Coded-Decimal

Decimal digit	BCD code
0	0000
1	0001
2	0010
3	0011
4	0100
5	0101
6	0110
7	0111
8	1000
9	1001

# Adição Usando BCD

$$\begin{array}{r} X \quad 0111 \quad 7 \\ + Y \quad + 0101 \quad + 5 \\ \hline Z \quad 1100 \quad 12 \\ \quad + 0110 \\ \hline \text{carry} \rightarrow 10010 \\ \quad \underbrace{\hspace{2em}} \\ \quad S = 2 \end{array}$$

$$\begin{array}{r} X \quad 1000 \quad 8 \\ + Y \quad + 1001 \quad + 9 \\ \hline Z \quad 10001 \quad 17 \\ \quad + 0110 \\ \hline \text{carry} \rightarrow 10111 \\ \quad \underbrace{\hspace{2em}} \\ \quad S = 7 \end{array}$$



# Adição Usando BCD

