# INSTITUTE OF COMPUTING - UNICAMP <br> Graduate Program 

MO417A Design and Analysis of Algorithms
2015 - Semester 1 - Jorge Stolfi
Problem Set 01-2015-03-09

| Nome | RA |  |  |  |  |  |  |  |  |  |  |  |  |
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| Item |  |  |  |  |  |  |  |  |  |  |  |  | TOT |
| Nota |  |  |  |  |  |  |  |  |  |  |  |  |  |

1. Consider the two sorting algorithms below, insertion sort $I$ and selection sort $S$. Both algorithms sort a list of $n$ numbers $x_{0}, x_{1}, \ldots, x_{n-1}$ in increasing order. The notation $[v]$ before a statement means " $v$ is the number of times that this statement gets executed when the algorithm is executed once".

| Algorithm $I(n, x)$ |  | Algorithm $S(n, x)$ |  |
| :---: | :---: | :---: | :---: |
| $\left[f_{I}\right]$ | for $i$ from 1 to $n-1$ do | [ $f_{S}$ ] | $i$ from 0 to |
|  | $t \leftarrow x_{i} ;$ |  | $k \leftarrow i ;$ |
|  | $j \leftarrow i$; |  | for $j$ from $i$ |
|  | while $j>0$ and $x_{j-1}>t$ do | [ $g_{S}$ ] | if $x_{j}<x^{\prime}$ |
| $\left[g_{I}\right]$ | $x_{j} \leftarrow x_{j-1} ;$ | $\left[h_{S}\right]$ | $k$ |
|  | $j \leftarrow j-1 ;$ |  | $\operatorname{swap} x_{i}, x_{k}$ |
|  | $x_{j} \leftarrow t ;$ |  |  |

(a) Write formulas for the counting variables $f_{I}, f_{S}, g_{I}, g_{S}$, and $h_{S}$, as a function of $n$, for If you cannot find exact formulas, or the value depends on the list $x$, find the best lower and upper bounds that you can.
$\square$
(b) What can you say about the relative efficiency of the two algorithms?
answer
2. This question refers to algorithm $H$ below, that takes as inputs a natural number $n$, and a list $X$ of $n$ integers, and returns an integer $m$. (The algorithm does not solve any interesting problem, it is just an exercise in analysis.)

```
    Algorithm \(H(n, X)\)
    \(m \leftarrow 0 ; i \leftarrow 0 ;\)
    while
[ \(f_{2}\) ]
        \(i<n\)
    do
        \(j \leftarrow i ;\)
        while
            \(j<n\)
        do
            \(A \leftarrow 0 ;\)
            \(r \leftarrow i ;\)
            while
[ \(\left.f_{7}\right] \quad r \leq j\)
            do
                \(A \leftarrow 2 A+X_{r} ;\)
                \(r \leftarrow r+1 ;\)
            endwhile;
[ \(f_{10}\) ] if \(A=0\) then
[ \(f_{11}\) ]
                \(m \leftarrow m+1 ;\)
            endif;
\(\left[f_{12}\right] \quad j \leftarrow j+1 ;\)
        endwhile;
[ \(f_{13}\) ] \(\quad i \leftarrow i+1 ;\)
    endwhile;
\(\left[f_{14}\right]\) return \(m\);
```

(a) Write the Kirchoff laws that relate the counting variables $f_{1}, f_{2}, \ldots$ (Drawing a flowchart of the algorithm may help.)
answer
(b) Identify a "key" subset of the counting variables $f_{1}, f_{2}, \ldots$, such that any other counting variable can be computed from that set, using the Kirchoff laws.
(answer
(c) Give formulas, as a function of $n$, for those key variables. If some variable depends on the matrix $X$, not only on $n$, give the best upper and lower bounds for that variable that you can find.
(1) answer

