Algorithm 1: Computation of orthonormal bases for each subspace V_* .

Data: Three $n \times n$ permutation matrices $A, B \in C$. **Result**: Subspace bases S_1, S_2, S_3, S_4, S_5 . $\mathbf{1} \ S \leftarrow \emptyset$ 2 $L \leftarrow \operatorname{null}(\begin{bmatrix} (A-B) \\ (A-C) \end{bmatrix})$ 3 foreach $v \in L$ do $\mathbf{4} \quad \left\lfloor \begin{array}{c} S \leftarrow \texttt{Add}(S, v) \end{array} \right.$ 5 $S_1 \leftarrow S //$ Basis of $V_*(.ABC.)$ 6 $S \leftarrow S_1$ 7 $L \leftarrow \operatorname{null}(A - B)$ 8 foreach $v \in L$ do $S \leftarrow \operatorname{Add}(S,v)$ 9 10 $S_2 \leftarrow S - S_1$ // Basis of $V_*(.AB.C.)$ **11** $S \leftarrow S_1$ 12 $L \leftarrow \operatorname{null}(B - C)$ 13 foreach $v \in L$ do 14 $S \leftarrow \operatorname{Add}(S,v)$ 15 $S_3 \leftarrow S - S_1$ // Basis of $V_*(.BC.A.)$ 16 $S \leftarrow S_1$ 17 $L \leftarrow \operatorname{null}(C - A)$ 18 foreach $v \in L$ do $S \leftarrow \operatorname{Add}(S,v)$ $\mathbf{19}$ 20 $S_4 \leftarrow S - S_1$ // Basis of $V_*(.AC.B.)$ **21** $L \leftarrow \operatorname{null}([S_1 \ S_2 \ S_3 \ S_4])$ **22** $S_5 \leftarrow \emptyset$ // Basis of $V_*(.A.B.C.)$ 23 foreach $v \in L$ do $S_5 \leftarrow \operatorname{Add}(S,v)$ $\mathbf{24}$

As v_2 is in the image of S_2 , it can be written as S_2u_2 , and therefore

$$P_2 v = S_2 (N_2^T S_2)^{-1} N_2^T S_2 u_2$$

= $S_2 u_2$
= v_2 .

Note that, in the computation of P_2 , P_3 , and P_4 , it is not possible to simplify the right-hand expression $S_i(N_i^T S_i)^{-1} N_i^T$. In general, neither S_i nor N_i are square matrices, and therefore they are not invertible. Nevertheless, the product $N_i^T S_i$ is always invertible.

Knowing the projection matrices, we can finally compute the median candidates, in Algorithm 2. Previously, we saw how to compute M_A . It is also possible to define M_B and M_C in an analogous way. The matrix M_B follows B in $V_*(.A.B.C.)$ instead of A, and M_C follows C. The entire computation takes $O(n^3)$ arithmetic operations.