

Shape-based Image Representation (Part I)

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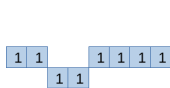
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- Each contour encodes important shape properties for object description.

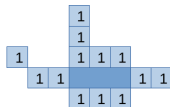
- Challenges in contour extraction and labeling.
- Algorithm for contour extraction and labeling.
- Geodesic length of a contour.

Challenges in contour extraction and labeling

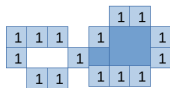
Noise in binary images might create (A) an open curve, (B) a contour with branches, and (C) a contour with bottleneck pixels. (D) Two contours might touch each other and (E) a contour might touch its opposite side.



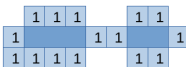
(A)



(B)



(C)



(D)



(E)

Challenges for contour pixel labeling

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- The dilation creates an image $\hat{J} = (D_I, J)$ such that

$$J(p) = \max_{q \in \mathcal{A}_r(p)} \{I(q)\},$$

where \mathcal{A}_r is from now on defined as

$$\mathcal{A}_r : \{(p, q) \in D_I \times D_I \mid \|q - p\| \leq r\}.$$

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- In any case, the algorithm must eliminate open curves, branches, ears at bottleneck pixels, and assign labels to pixels of touching contours and of contours that touches their opposite side.

Contour pixel labeling

- For a given binary image $\hat{I} = (D_I, I)$, the algorithm must be constrained to the **border set** \mathcal{B} .

$$\mathcal{B} : \{p \in D_I \mid I(p) = 1 \text{ and } \exists q \in \mathcal{A}_1(p) \mid I(q) = 0\}.$$

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- For each contour, the algorithm must return a path $\pi_{p_1 \rightarrow p_n}$, wherein $(p_i, p_{i+1}) \in \mathcal{A}_{\sqrt{2}}$, $i = 1, 2, \dots, n - 1$, being n the size of the contour.

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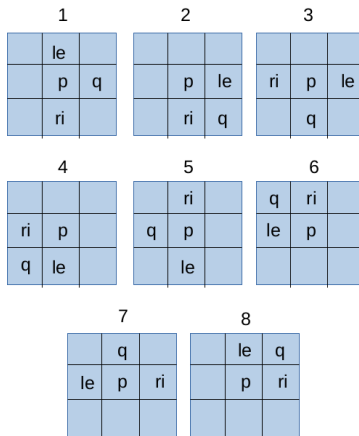
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- It must also avoid arcs $\langle p, q \rangle$ which pass through the object/background and take into account the anti-clockwise orientation (i.e., $I(le(p, q)) = 1$ and $I(ri(p, q)) = 0$).

Contour pixel labeling

It starts from pixels that define valid arcs and visits the border pixels in **depth search** using **clockwise** adjacency relation $\mathcal{A}_{\sqrt{2}}$.

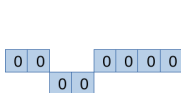
Clockwise adjacency relation



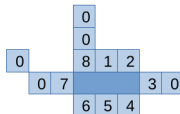
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Pixel labeling result



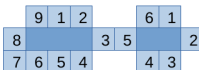
(A)



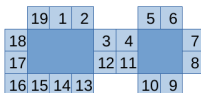
(B)



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The first two lines of the algorithm are dedicated to define the border set \mathcal{B} (**initially empty**), initialize a predecessor map P , and a label map L_{px} .

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1. For each $p \in D_I$, set $P(p) \leftarrow nil$ and $L_{px}(p) \leftarrow 0$.
2. If $I(p) = 1$ and $\exists q \in \mathcal{A}_1(p) \mid I(q) = 0$, insert p in \mathcal{B} .

Contour pixel labeling

- Two conditions are used for a border pixel p and an arc $\langle p, q \rangle$ be considered a **valid starting pixel** and a **valid arc**, respectively.
 - Valid starting pixel p : $p \in \mathcal{B}$, $P(p) = nil$, and $\exists q \in \mathcal{A}_{\sqrt{2}}(p) \mid \langle p, q \rangle$ is a valid arc and q could also be a starting pixel.

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 - Valid arc $\langle p, q \rangle$: $q \in \mathcal{A}_{\sqrt{2}}(p)$, $q \in \mathcal{B}$, $P(q) = nil$, $I(le(p, q)) = 1$, and $I(ri(p, q)) = 0$.

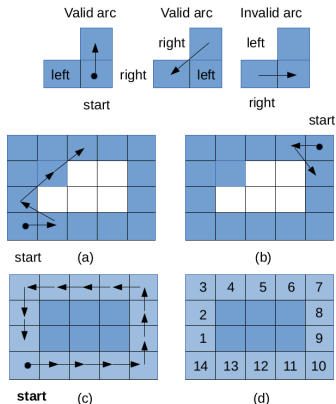
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 - Valid arc $\langle p, q \rangle$: $q \in \mathcal{A}_{\sqrt{2}}(p)$, $q \in \mathcal{B}$, $P(q) = nil$, $I(lc(p, q)) = 1$, and $I(ri(p, q)) = 0$.
- Let's indicate these events by $Valid(p)$ and $Valid(p, q)$.
- Such rules cannot prevent the algorithm to start labeling silly contours, such as $(x, y) \rightarrow (x + 1, y) \rightarrow (x, y - 1)$, $(x, y) \rightarrow (x, y - 1) \rightarrow (x - 1, y)$, etc.

Contour pixel labeling



Silly contour with a last invalid arc at the top. Labels should not be assigned when (a) the path starts but never finishes and (b) it finishes as a silly contour. (c)-(d) A valid contour and its pixel labels.

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- At this moment, if it is not a silly contour, we can start the labeling process.
- A silly contour is detected when $P(P(p))$ is the starting pixel.
- Therefore, the algorithm visits pixels in \mathcal{B} in anti-clockwise until the starting pixel is reached again. If the path is not a silly contour, its pixels are labeled clockwise using P .

Contour pixel labeling

03. For each $s \in D_I \mid \text{Valid}(s)$, do.
04. Set $P(s) \leftarrow s$ and insert s in Q .
05. While $Q \neq \emptyset$, do.
06. Remove p from Q .
07. For each $q \in \mathcal{A}_{\sqrt{2}}(p)$, do.
08. If $q = s$ and $P(p) \neq s$, then
09. If $P(P(p)) \neq s$, then go to 13, else go to 17.
10. If $\text{Valid}(p, q)$, then.
11. Set $P(q) \leftarrow p$.
12. Insert q in Q .
13. Set $i \leftarrow 1$.
14. While $P(p) \neq p$, do.
15. Set $L_{px}(p) \leftarrow i$, $p \leftarrow P(p)$ and $i \leftarrow i + 1$.
16. Set $L_{px}(p) \leftarrow i$.
17. Set $Q \leftarrow \emptyset$.

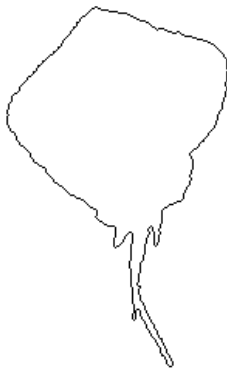
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One can easily change it to assign a distinct contour label.

Contour pixel labeling

An example of the shape of a fish (left) and its contour pixels labeled from 1 to n in clockwise orientation (right) — brighter the pixel lower is the label.



Geodesic length of a contour

- Let (p_n, p_1) be a valid arc in a border set \mathcal{B} and $\Pi_{\mathcal{B}}$ be the set of all possible paths in a graph $(\mathcal{B}, \mathcal{A}_{\sqrt{2}})$.

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- The length $f_{geo}(\pi_{p_1 \rightsquigarrow p_n}^*)$ of $\pi_{p_1 \rightsquigarrow p_n}^*$ is called **geodesic** [1].
- The geodesic length of a contour in \mathcal{B} is defined as $f_{geo}(\pi_{p_1 \rightsquigarrow p_n}^*) + w(p_n, p_1)$, where

$$f_{geo}(\pi_{p_1 \rightsquigarrow p_n}^*) = \min_{\pi_{p_1 \rightsquigarrow p_n} \in \Pi_{\mathcal{B}} \setminus \langle p_1, p_n \rangle} \left\{ \sum_{k=1}^{n-1} w(p_k, p_{k+1}) \right\},$$
$$w(p_k, p_{k+1}) = \begin{cases} 0.9016 & \text{if } \|p_{k+1} - p_k\| = 1, \\ 1.2890 & \text{if } \|p_{k+1} - p_k\| = \sqrt{2}. \end{cases}$$

Geodesic length of a contour

- The shorest path formulation can label each pixel p_k , $k \in [1, n]$, of a contour $\pi_{p_1 \rightsquigarrow p_n}^* \cdot \langle p_n, p_1 \rangle$ by the geodesic length $f_{geo}(\pi_{p_1 \rightsquigarrow p_k}^*)$.

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- We will see how useful is this formulation to obtain smooth multiscale skeletons from the shape.
- As exercise, elaborate an IFT-based algorithm to extract and label all contours in a binary image by the geodesic length assigned to each pixel with respect to an arbitrary pixel p_1 and a valid arc $(p_n, p_1) \in \mathcal{A}_{\sqrt{2}}$.

[1] A.X. Falcão, C. Feng, J. Kustra, and A.C. Telea.

Chapter 2 - multiscale 2d medial axes and 3d surface skeletons by the image foresting transform.

In P.K. Saha, G. Borgefors, and G.S. di Baja, editors, *Skeletonization*, pages 43 – 70. Academic Press, 2017.