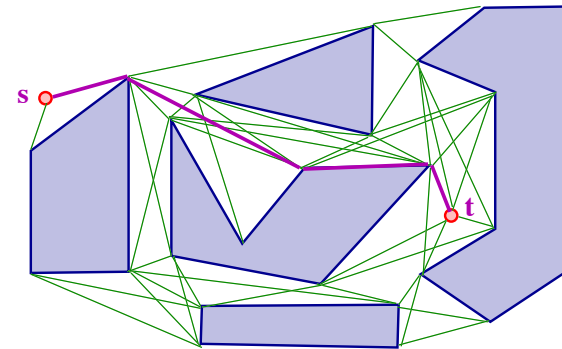


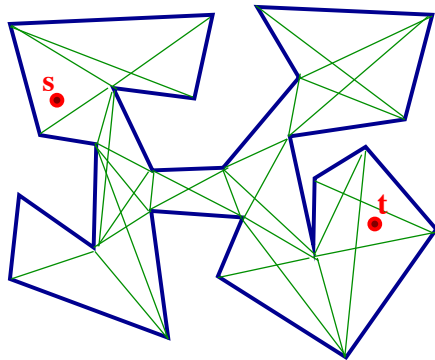
Visibility Graph



- Construct a visibility graph $G = (V, E)$, where V is set of polygon vertices (and s, t), E is pairs of nodes that are mutually “visible”.
- Give each edge (u, v) the weight equal to the Euclidean distance between u and v .
- The shortest path from s to t in this graph is the obstacle avoiding shortest path.
- G can have between c_1n and c_2n^2 edges. Run Dijkstra’s algorithm.

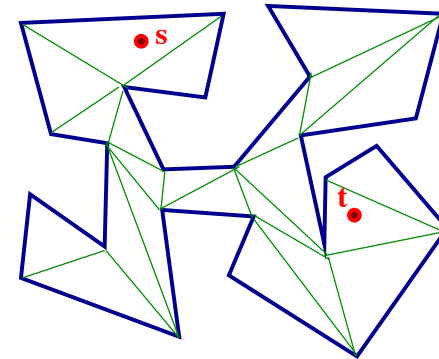
Paths in a Polygon

- Workspace interior of a simple polygon.
- Can we compute a shortest path faster?
- The visibility graph can still have $\Theta(n^2)$ edges.



- Using polygon triangulation, we show an $O(n \log n)$ time algorithm.

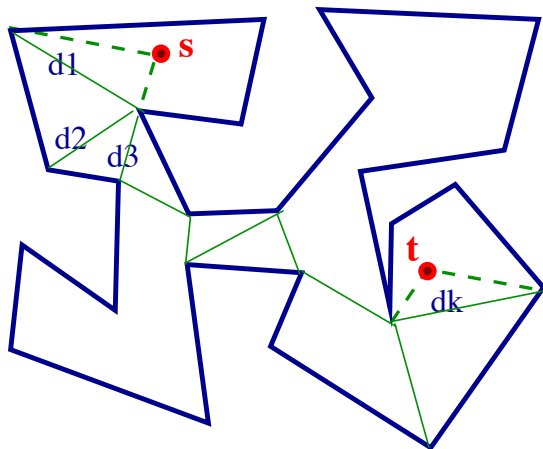
Fast Algorithm



- Let P be a simple polygon and s, t be source and target points.
- Let T be a triangulation of P .
- Call a **diagonal d** of T **essential** if s, t lie on opposite sides of d .
- Let d_1, d_2, \dots, d_k be ordered list of essential diagonal.

Algorithm

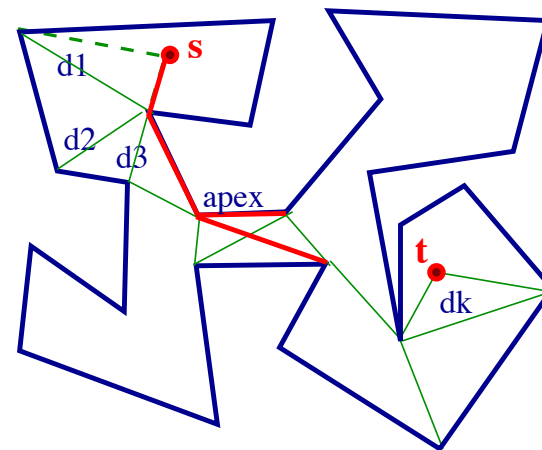
- Essential diagonals d_1, d_2, \dots, d_k .



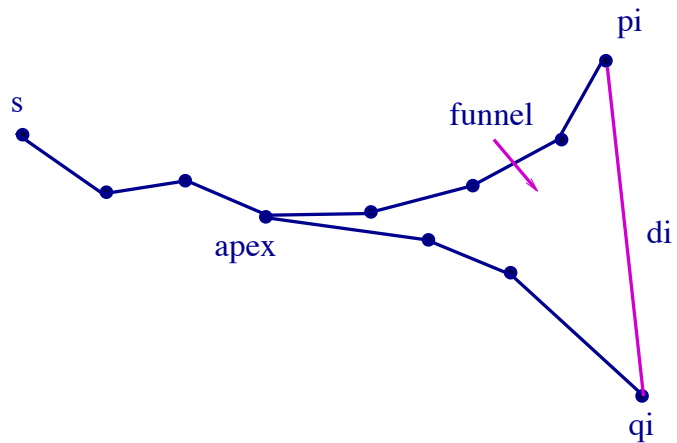
- The algorithm works as follows:
 1. Start with $d_0 = s$.
 2. for $i = 1$ to $k + 1$ do
 3. Extend path from s to both endpoints of d_i

Path Extending: Funnel

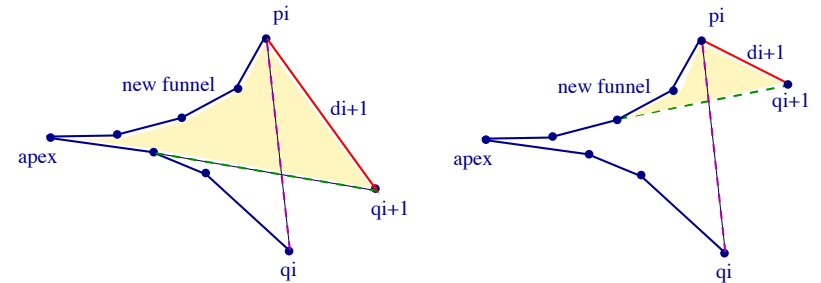
- Union of $path(s, p_i)$ and $path(s, q_i)$ forms a **funnel**.
- The vertex where paths diverge is called **apex**.



Funnel

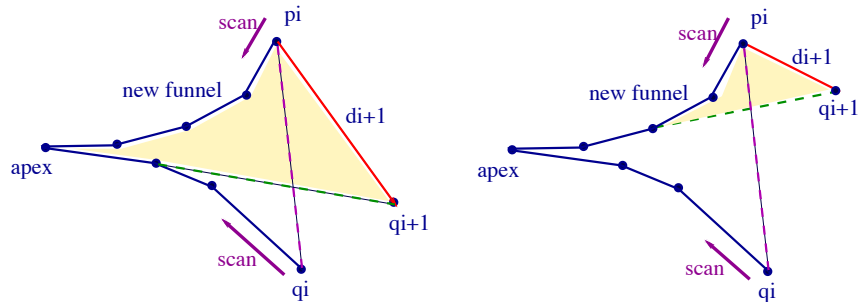


Path Extending



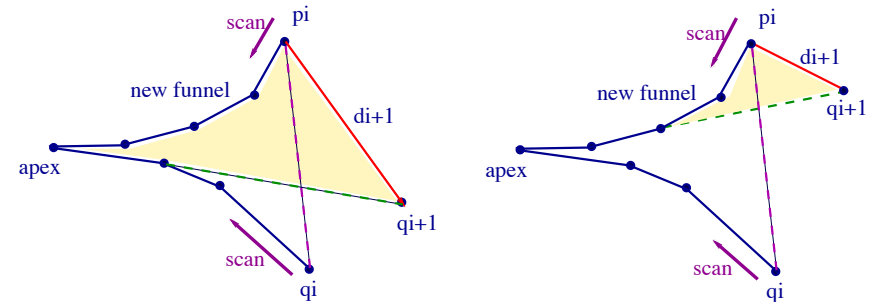
- Two cases of how to extend the path.
- In case I, funnel contracts.
- In case II, apex shifts, tail extends, funnel contracts.
- In each case, funnel property maintained.

Data Structure & Update



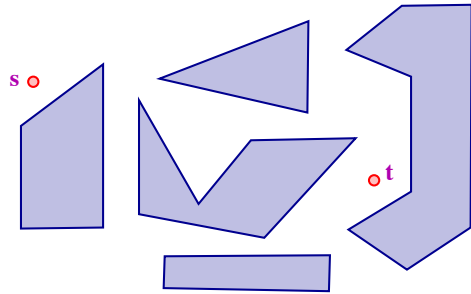
- How to determine tangent to funnel?
- Can't afford to spend $O(n)$ time for each tangent.
- Idea: If x edges of funnel are removed by the new tangent, spend $O(x)$ time for finding the tangent.
- How to tell a tangent?

Data Structure & Update



- Start scanning the funnel from both ends, until tangent determined.
- At most $2x + 2$ vertices scanned.
- Since each vertex inserted once, and deleted once, total cost for all the tangents is $O(n)$.
- Data structure for the funnel: Double-ended queue. Insert/delete in $O(1)$ time.

Paths Among Obstacles



Approach	Complexity	Reference
Vis. Graph	$O(n^3)$ $O(n^2 \log n)$ $O(n^2)$ $O(E + n \log n)$	L. Perez, Wesley '79 Sharir-Schorr '84 Welzl, AAGHI '86 Ghosh-Mount '91
SP Map	$O(k^2 + n \log n)$ $O(nk \log n)$ $O(n^{5/3+\epsilon})$ $O(n \log n)$	Kapoor-Maheshwar Reif-Storer '91 Mitchell '93 Hershberger-S '93